

(12) INTERNATIONAL APPLICATION PUBLISHED UNDER THE PATENT COOPERATION TREATY (PCT)

(19) World Intellectual Property Organization
International Bureau(43) International Publication Date
19 July 2001 (19.07.2001)

PCT

(10) International Publication Number
WO 01/52496 A2

- (51) International Patent Classification⁷: H04L 29/00
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- (21) International Application Number: PCT/US00/33640
- (81) Designated States (national): AE, AL, AM, AT, AU, AZ, BA, BB, BG, BR, BY, CA, CH, CN, CR, CU, CZ, DE, DK, DM, EE, ES, FI, GB, GD, GE, GH, GM, HR, HU, ID, IL, IN, IS, JP, KE, KG, KP, KR, KZ, LC, LK, LR, LS, LT, LU, LV, MA, MD, MG, MK, MN, MW, MX, NO, NZ, PL, PT, RO, RU, SD, SE, SG, SI, SK, SL, TJ, TM, TR, TT, TZ, UA, UG, UZ, VN, YU, ZA, ZW.
- (22) International Filing Date:
11 December 2000 (11.12.2000)
- (25) Filing Language: English
- (26) Publication Language: English
- (30) Priority Data:
09/479,781 7 January 2000 (07.01.2000) US
- (84) Designated States (regional): ARIPO patent (GH, GM, KE, LS, MW, MZ, SD, SL, SZ, TZ, UG, ZW), Eurasian patent (AM, AZ, BY, KG, KZ, MD, RU, TJ, TM), European patent (AT, BE, CH, CY, DE, DK, ES, FI, FR, GB, GR, IE, IT, LU, MC, NL, PT, SE, TR), OAPI patent (BF, BJ, CG, CI, CM, GA, GN, GW, ML, MR, NE, SN, TD, TG).
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- Published:**
— without international search report and to be republished upon receipt of that report
- For two-letter codes and other abbreviations, refer to the "Guidance Notes on Codes and Abbreviations" appearing at the beginning of each regular issue of the PCT Gazette.*



WO 01/52496 A2

(54) Title: A DECLARATIVE LANGUAGE FOR SPECIFYING A SECURITY POLICY

(57) Abstract: The invention is a declarative language system and comprises a language as a tool for expressing network security policy in a formalized way. It allows the specification of security policy across a wide variety of networking layers and protocols. Using the language, a security administrator assigns a disposition to each and every network event that can occur in a data communications network. The event's disposition determines whether the event is allowed (i.e. conforms to the specified policy) or disallowed and what action, if any, should be taken by a system monitor in response to that event. Possible actions include, for example, logging the information into a database, notifying a human operator, and disrupting the offending network traffic.

A DECLARATIVE LANGUAGE FOR SPECIFYING A SECURITY POLICY

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BACKGROUND OF THE INVENTION

TECHNICAL FIELD

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The invention relates to security and network services. More particularly, the invention relates to a declarative language system used in defining policy for an entire network and in providing monitoring and enforcing of computer network security.

15

DESCRIPTION OF THE PRIOR ART

Security administrators need tools that help them formulate their site security policy and translate it into monitoring and enforcement mechanisms. They
20 need to be sure that the computer enforced policy – often cobbled together from a plethora of disjoint access control mechanisms – matches their enterprise policy, all too often specified in a loose natural language or a set of unwritten principles. This leads to confusion as to why access is being granted or denied to particular resources and may lead to unintentional
25 breaches of security.

A way to reduce or eliminate the confusion described above is by providing a user-friendly and, yet, rigorous way of specifying security policy, as well as providing tools for monitoring and enforcing the security policy.

30

Blaze, Feigenbaum, and Lacy (BFL), *Decentralized Trust Management*, Proc. IEEE Conference on Security and Privacy (1996), used the term *trust management* to refer to a problem of deciding whether requested actions, supported by credentials, conform to policies. In other words, it deals with the

questions of who, how, and what. *Who* (the principals, for example, people, computers and organizations) can access *what* (the resources being sought) and *how* (the actions performed against the target resources).

- 5 Mansouri-Samani, *et al*, *GEM: A Generalized Monitoring Language for Distributed Systems*, Distributed Systems Engineering, vol.4, no. 2 96-108 (June 1997) discloses a generalized-event monitoring notation that permits user-specified filtering and composition scripts to be dynamically loaded into distributed-event monitoring components. GEM uses "scheduled time events
10 and default or user-defined detection windows" to cope with "variable communication delay problems." The GEM event monitoring system is used "to detect complex event sequences and to convert these into simple events" that trigger management actions. The event monitors have been restricted to performing "very simple activities related to triggering or notifying events."

- 15 J. A. Grompone, *A Declarative Language for the Configuration of Exchanges*, Telecommunications Journal, vol. 56, no.1 (Jan. 1989) discloses the design and implementation of a high-level language, LEP, to define the routing and customizing of rules of a telex exchange. The routing concepts are basic and
20 few in number. Each of the physical communication paths is called a line. The lines are arranged in groups. The purpose of the LEP language is to provide a comprehensive definition of all lines of an exchange, the arrangement of these lines in groups and the physical attributes of the groups. All groups taken together comprise all the lines without any lines missing or
25 being repeated. A group is an ordered set of lines. The LEP term "access" is used to denote whether lines are permitted or forbidden to access other lines or services. Routing, a basic objective of an LEP program, is a way of associating sets of compiled codes with destinations, done through a sequence of elementary declarations. LEP also defines the possible
30 destinations of a call. One of the main design concepts was to use a very simple structure for the declarations for even users unfamiliar with computer programming.

- 35 The LEP language cannot thread together multiple protocol layers of a network event. The LEP language lacks the sophistication in terms of richer

expressions to allow a set of policy rules affecting different networking protocols to be applied to a complex protocol interaction between two communicating parties, and to security policy for an entire network. The LEP language does not suggest defining allowed traffic patterns and handling those events that deviate from those patterns.

Plasek, *et al*, *Statistical Database Query Using Random Sampling Of Records*, U.S. Patent 5,878,426, discloses a method for obtaining decision support query results from a database table having multiple records. An attribute of the database table is sampled, which results in a collection of sampled data. The sampled data represents some percentage of all of the data corresponding to that attribute in the database table. The data associated with the attribute includes multiple data classes, and the sampled data is separated or partitioned into these data classes. A database query is applied to the sampled data rather than to all of the data corresponding to that attribute in the database table.

Plasek, *et al*, also discloses a method to obtain decision support query results from a database table where all of the data associated with a particular database attribute is grouped into various data classes. Each of the data classes is individually randomly sampled to obtain a corresponding number of class data samples. Each of the class data samples is then queried, which can include executing aggregation functions on each of the class data samples.

Plasek, *et al*, also discloses a method for providing result approximations in database queries.

Plasek, *et al*, does not disclose nor suggest providing a method to select a most specific and applicable result or policy rule. Plasek, *et al*, does not disclose nor suggest providing a method to rank data and does not order data in a database beyond partitioning data into classes and thereafter randomly sampling each data class such that database queries are applied to each of the samples.

Plasek, *et al*, does not disclose nor suggest providing a method to thread protocol layers of a network event together to provide a result to the network event.

- 5 Chow, *et al*, *System, Method, and Program for Extending a SQL Compiler for Handling Control Statements Packaged with SQL Query Statements*, U.S. Patent No. 5,875,334 (February 23, 1999) discloses an integrated compiler for compiling SQL3 control statements having procedural, i.e., control, information packaged together with query, i.e., non-procedural, statements. A
10 query extractor contained within the parser extracts the query statement from the control statement leaving a control skeleton. The query statement is processed as usual through a query compiler for generating executable plans with the exception that the name resolution function for resolving variables is modified for looking up local variables. This modification takes into account
15 the mapping of local and host variables to create a unification of local and host variables. The control skeleton is processed through a control analyzer which generates a representation of the control flow and a scope and symbol table. The control analyzer also unifies the local and host variables. A plan synthesizer then takes as input the control flow information, symbol tables, and individual executable plans for the query statements and generates a
20 meta-plan comprising a merger of a top level plan for the control skeleton and sub-plans representing the executable plans of the query statement.

- Chow, *et al*, does not disclose nor suggest a ranking method or an ordering
25 method to handle a set of rules to be applied to a complex protocol interaction between two communicating parties.

- Nor does Chow, *et al*, disclose or suggest a method whereby to thread protocol layers of a network event together to provide a rule applicable to the
30 network event.

- V. Paxson, *Bro: A System for Detecting Network Intruders in Real-Time*, Network Research Group, Lawrence Berkeley National Laboratory, Berkeley, CA, LBNL-41197 (Jan. 1998) discloses a stand-alone system for detecting
35 network intruders in real-time by passively monitoring a network link over

which the intruder's traffic transits. The system comprises a "policy script interpreter" that interprets event handlers written in a specialized language used to express a site's security policy. The specialized language is C-style because it comprises, for example, C-style data types and constants, operators, and block statements and is procedural. Bro comprises first-class values and aggregate types such as record and table, used to specify a security policy.

However, Paxson does not disclose nor suggest providing a sophisticated ranking method to rank policy rules according to the specificity of the initiator and target communicating hosts and to select a most applicable rule in an efficient manner. Paxson does not disclose nor suggest providing a method to thread protocol layers of a network event together to provide a result to the entire network event.

It would be advantageous to reduce or eliminate the confusion described herein above by providing a user-friendly and, yet, rigorous way of specifying security policy, as well as providing tools for monitoring and enforcing the security policy.

It would be advantageous to have a trust manager that takes as its input a security policy defined as a set of *policy rules* (statements about trust) and a set of *credentials* (statements about principals), such that it is capable of processing requests for *trust decisions*, i.e. evaluating compliance with the policy.

It would be advantageous for the trust manager to have a unified view of an interaction between two principals across a stack of protocol layers, each governed by discreet policy rules, and to apply a final trust decision based on which of these policy rules better fits the entire interaction. For example, using HTTPS to access a secure web page involves an interaction between two network addressable machines (at the TCP/IP level), an interaction between a cryptographically authenticated server and, possibly, a cryptographically authenticated client (at the SSL level), and an interaction

between a Web browser (possibly with its own authentication credentials) and a web server, resulting in the retrieval of a web page (at the HTTP level).

5 It would be advantageous to have a *policy definition language* as well as well-designed algorithms to support monitoring and auditing network activity, in addition to traditional access/deny authorization decisions. For example, a policy rule might instruct a monitoring Agent to log all traffic between two computers or to decrypt a secure channel between two users.

10 It would be advantageous to provide a system that comprises a passive monitor of network traffic that does not need to be installed on target hosts or integrated into existing applications.

15 It would be advantageous to provide a system that uses a sophisticated algorithm for determining which policy rules take precedence over others.

20 It would be advantageous to provide a policy language that allows a set of policy rules affecting different networking protocols to be applied to a complex protocol interaction between two communicating parties. Also, it would be advantageous to use the policy language to express security policy for an entire network.

25 It would be advantageous to provide a system, unlike current Intrusion Detection Systems (IDS) which only look for signatures of known attacks, focusing on defining allowed traffic patterns and determining how to handle events that deviate from those patterns.

SUMMARY OF THE INVENTION

30 The invention is a declarative language system and comprises a language as a tool for expressing network security policy in a formalized way. It allows the specification of security policy across a wide variety of networking layers and protocols. Using the language, a security administrator assigns a disposition to each and every network event that can occur in a data communications

network. The event's disposition determines whether the event is allowed (*i.e.* conforms to the specified policy) or disallowed and what action, if any, should be taken by a system monitor in response to that event. Possible actions include, for example, logging the information into a database, notifying a human operator, and disrupting the offending network traffic.

The language is implemented by a Policy Engine, a component also of a Security Policy Monitoring (SPM) system. The SPM system, also referred to herein as the Policy Monitoring System, is ideally suited for network and security assessments where real network traffic is analyzed in order to identify abnormal traffic patterns, system vulnerabilities and incorrect configuration of computer systems on the network.

Unlike other trust management systems, the SPM is designed to be a passive monitor of network traffic. As such, it need not be installed on target hosts or integrated into existing applications.

The invention provides a simple and intuitive model for expressing and applying security policies. The language is much richer in terms of what it can express than languages used in firewalls and routers. It uses a sophisticated algorithm for determining which policy rules take precedence over others, a process that in other systems is completely manual.

Unlike existing firewalls and routers, the invention allows a set of policy rules affecting different networking protocols to be applied as a whole to a complex protocol interaction between two communicating parties. Furthermore, networking equipment typically handles only policy related to the network traffic that flows through it. Using the invention herein one can express the security policy for an entire network.

Unlike IDS systems, which look for the signatures of known attacks, the SPM is focused on defining allowed traffic patterns and how to handle events that deviate from those patterns.

BRIEF DESCRIPTION OF THE DRAWINGS

Fig. 1 is a schematic diagram showing the relationship of elements of the Policy Monitoring System, according to the invention;

5

Fig. 2 is a schematic diagram of a protocol event according to the invention;

Fig. 3 is a schematic diagram of a disposition according to the invention;

10 Fig. 4 is a schematic diagram of communicating parties according to the invention;

Fig. 5a is a schematic diagram of a network event, comprising protocol events at different protocol layers, having an associated network event disposition according to the invention; and

15

Fig. 5b is an algorithm showing protocol events at different protocol layers resulting in pending rules with or without immediate outcomes and, finally, a final disposition for the network event.

20

DETAILED DESCRIPTION OF THE INVENTION**Overview**

25 Fig. 1 is a schematic diagram showing the relationship of elements of the Policy Monitoring System 100, according to the preferred embodiment of the invention. To effect a security policy decision, a policy manager module invokes a Policy Engine 101 with both a reference to a pre-defined security policy and a number of inputs it received from an Agent 102. These inputs

30 describe a protocol event such as a TCP connection, an SSL session, or an HTTP GET. The end to end interaction between two communicating entities comprises one or more protocol events and it is termed a network event 103. For example, the retrieval of a web page from a web server by a web browser

is a network event that typically consists of an IP protocol event, a TCP protocol event and an HTTP protocol event.

5 In the preferred embodiment the Policy Engine 101 consults a policy information database, a Policy Store 104 to determine a policy rule that applies to the network event 103. In the preferred embodiment the Policy Engine 101 collects input from the Agent 102 about each protocol event until it has enough information to consult the Policy Store 104. Once an applicable policy rule for the entire network event 103 has been found, the Policy Engine
10 101 returns a disposition 105 for the event to the policy manager module which in turn forwards it to the Agent 102, to a logging subsystem and, optionally, to an enforcement subsystem.

A definition of a protocol event is provided to facilitate understanding of the
15 invention. A protocol event 120 as shown in Fig. 2 comprises the following elements:

1) The Principals. Every policy decision involves two principals: an initiator 121 (an active principal) and a target 122 (a passive principal). Principals are
20 identified by a set of one or more credentials, depending on how much information is known about them and what protocol service they are using. There are three types of credentials:

- a) Host credentials. These are the physical network address, *i.e.* a MAC address, and the network attachment point, *e.g.* an IP address and port
25 number.
- b) User credentials. These may be weak credentials, *e.g.* a user name and password, or strong credentials, *e.g.* an X.509 certificate.
- c) Object credentials. These identify a resource or an application, *e.g.* a URL or a pathname.

30 2) The Protocol. The protocol service 123 associated with this protocol event 120.

3) The Security Quality of Service Parameters. Some protocols include
35 security QOS parameters 124 and these may be subject to local security

policy constraints. For example, in the SSL protocol the ciphersuite negotiated between the SSL client and the SSL server is a security QOS parameter.

- 5 4) The Action. Every interaction between an initiator and a target over a given protocol service involves a specific action 125. Clearly, not all actions are of interest to the policy manager module. For example, in the SSL protocol only actions pertaining to the establishment or termination of an SSL session, most notably, the negotiation of security parameters for the session
10 are of interest. In the LDAP protocol, on the other hand, a security policy administrator may wish to express policy statements about different LDAP data manipulation operations, such as, the SEARCH and MODIFY operations.

- In one embodiment of the invention, while processing a network event 103,
15 and before issuing a final ruling, the Policy Engine 101 may instruct the Agent 102 to carry out specific actions against the network event 103. For example, the Agent 102 may be asked to decrypt subsequent SSL traffic or it may be asked to impose a specific ciphersuite on the target system. These instructions constitute an intermediate output of the Policy Engine 101 and are
20 issued in the form of agent directives, defined herein below.

Once the Policy Engine 101 arrives at a final policy decision, it produces a disposition 105 for the event 103. The disposition 105 as shown in Fig. 3 comprises the following elements:

25

- 1) Disposition Code. The disposition code 131 denotes whether or not the event 103 complies with the security policy and, if not, identifies the specific policy violation. A list of possible codes in a preferred embodiment is given in Table A herein below. This field is mandatory.

30

- 2) Logging Directives. The logging directives field 132 includes a severity code, denoting the severity of the policy violation. A list of possible severity values in a preferred embodiment is given herein below in Table B. The severity code may be used by a logging subsystem to filter the event 103 and
35 its disposition 105 or to control a notification action, e.g. page a network

operator. In another embodiment the logging directives 132 may also include an optional human readable description summarizing the specific policy that determined the final disposition 105 e.g. "blue users cannot access the red server". The logging directives field 132 is mandatory if the disposition code

5 131 indicates a policy violation.

3) Agent Directives. Agent directives 102 are any instructions that need to be communicated to the Agent 102 and in another embodiment to more than one Agent. For example, the Agent 133 may be instructed to log all traffic associated with the event 103 or to disrupt communications between the initiator 121 and the target 122. In some embodiments, an Agent 102 only supports monitoring functions, or only enforcement functions, or be limited in its support of other types of functions. In a preferred embodiment, a policy manager is responsible for distributing a set of directives to appropriate Agents.

4) Event Constraints. Event constraints 134 are any constraints to be applied to the event 103. For example, in one embodiment these constraints are protocol-specific constraints such as the maximum lifetime of a TCP connection or the maximum lifetime of an SSL session. In another embodiment, these constraints are communicated to the Agent reporting the event or simply to a policy manager.

Table A

25 Built-in Objects

The following is a set of built-in language objects known to both the policy compiler and the policy engine.

First the built-in groups. It should be noted that, unlike user-defined groups, built-in groups cannot be extended in a policy specification.

```
// List of supported protocols
( group all-protocols protocol_t
  ( union IP UDP ICMP TCP SSL HTTP )
  // NOTE: new protocols can be added as needed
)
```

```
// List of supported hash algorithms
( group hash-algorithms hash_alg_t
  ( union MD5 SHA1 )
5 )

// List of supported agent directives
( group agent-directives agent_directive_t
  ( union DECRYPT DISRUPT LOG_TRAFFIC )
10 )

// List of supported logging severity codes
( group severity-codes severity_t
  ( union CRITICAL HIGH MEDIUM WARNING MONITOR INFORMATION )
15 )

// List of supported disposition codes
( group disposition-codes code_t
  ( union OK CONTINUE ACCESS_DENIED AUTHENTICATION_VIOLATION
20 SECURITY_ATTACK SECURITY_QOS POLICY_ERROR )
  )

// Certificate status values for valid certificates
( group valid-certs cert_status_t
25 ( union VALID )
  )

// Certificate status values for certificates rendered invalid
( group invalid-certs cert_status_t
30 ( union EXPIRED NOT_YET_VALID REVOKED SUSPENDED )
  )

// Certificate status values for rejected certificates
( group rejected-certs cert_status_t
35 ( union MALFORMED UNSUPPORTED_CRITICAL_EXTENSION )
  )

// Certificate status values for all bad certificates
( group bad-certs cert_status_t
40 ( union rejected-certs invalid-certs )
  )
```

```

// List of all possible certificate status values
( group cert-status-values cert_status_t
    ( union valid-certs invalid-certs rejected-certs )
5 )

// List of all possible authentication status values
( group auth-status-values auth_status_t
    ( union SUCCEEDED REJECTED ABORTED )
10 )

// List of all SSL ciphersuites
( group ssl-ciphersuites cipher_suite_t
    ( union
15     SSL_RSA_WITH_NULL_MD5
        SSL_RSA_WITH_NULL_SHA
        SSL_RSA_EXPORT_WITH_RC4_40_MD5
        SSL_RSA_EXPORT_WITH_RC2_CBC_40_MD5
        SSL_RSA_EXPORT_WITH_DES40_CBC_SHA
        SSL_RSA_WITH_RC4_128_MD5
20     SSL_RSA_WITH_RC4_128_SHA
        SSL_RSA_WITH_IDEA_CBC_SHA
        SSL_RSA_WITH_DES_CBC_SHA
        SSL_RSA_WITH_3DES_EDE_CBC_SHA
        SSL_DH_RSA_WITH_3DES_EDE_CBC_SHA
25     SSL_DH_DSS_WITH_3DES_EDE_CBC_SHA
        SSL_DH_RSA_WITH_DES_CBC_SHA
        SSL_DH_DSS_WITH_DES_CBC_SHA
        SSL_DH_RSA_EXPORT_WITH_DES40_CBC_SHA
        SSL_DH_DSS_EXPORT_WITH_DES40_CBC_SHA
30     SSL_DH_ANON_EXPORT_WITH_RC4_40_MD5
        SSL_DH_ANON_WITH_RC4_128_MD5
        SSL_DH_ANON_EXPORT_WITH_DES40_CBC_SHA
        SSL_DH_ANON_WITH_DES_CBC_SHA
        SSL_DH_ANON_WITH_3DES_EDE_CBC_SHA
35     SSL_DHE_RSA_WITH_3DES_EDE_CBC_SHA
        SSL_DHE_DSS_WITH_3DES_EDE_CBC_SHA
        SSL_DHE_RSA_EXPORT_WITH_DES40_CBC_SHA
        SSL_DHE_DSS_EXPORT_WITH_DES40_CBC_SHA
        SSL_DHE_RSA_WITH_DES_CBC_SHA
40     SSL_DHE_DSS_WITH_DES_CBC_SHA
        SSL_FORTEZZA_KEA_WITH_NULL_SHA
    )
    )

```

```

        SSL_FORTEZZA_KEA_WITH_FORTEZZA_CBC_SHA
        SSL_FORTEZZA_KEA_WITH_RC4_128_SHA
        SSL_V2_RC4_128_WITH_MD5
        SSL_V2_RC4_128_EXPORT40_WITH_MD5
5      SSL_V2_RC2_CBC_128_CBC_WITH_MD5
        SSL_V2_RC2_CBC_128_CBC_EXPORT40_WITH_MD5
        SSL_V2_IDEA_128_CBC_WITH_MD5
        SSL_V2_DES_64_CBC_WITH_MD5
        SSL_V2_DES_192_EDE3_CBC_WITH_MD5
10    )
    )

    // List of supported action codes for TCP
    ( group tcp-action-codes action_t
15      ( union CONNECT
          MISSED_CONNECT
          TIMEOUT
          ABORT
          CLOSE )
20    )

    // List of supported action codes for UDP
    ( group udp-action-codes action_t
        ASSOCIATION
25    )

    // List of supported action codes for IP
    ( group ip-action-codes action_t
        ASSOCIATION
30    )

    // List of supported action codes for ICMP
    ( group icmp-action-codes action_t
        ( union ASSOCIATION
35          BAD_CODE
          FRAGMENTATION_NEEDED
          HOST_UNREACHABLE
          NETWORK_UNREACHABLE
          PORT_UNREACHABLE
40          PROTOCOL_UNREACHABLE
          SOURCE_ROUTE_FAILED

```

```

ECHO
ECHO_REPLY
INFORMATION_REQUEST
INFORMATION_REPLY
5   PARAMETER_PROBLEM
    REDIRECT_HOST
    REDIRECT_TYPE_OF_SERVICE_AND_HOST
    REDIRECT_NETWORK
    REDIRECT_TYPE_OF_SERVICE_AND_NETWORK
10  SOURCE_QUENCH
    TIME_TO_LIVE_EXCEEDED
    REASSEMBLY_TIME_EXCEEDED
    TIMESTAMP
    TIMESTAMP_REPLY )
15 )

    // List of supported action codes for SSL
    ( group ssl-action-codes action_t
      ( union HANDSHAKE
20      MISSED_HANDSHAKE
        SESSION_CLOSED
        SESSION_ABORTED )
      )

25 // List of supported action codes for HTTP
    ( group http-action-codes action_t
      ( union GET
        HEAD
        POST
30      PUT
        DELETE
        OPTIONS
        TRACE
        CONNECT
35      MISSED_REQUEST
        RESPONSE )
      )

    // List of all supported action codes
40 ( group all-action-codes action_t
      ( union udp-action-codes

```



```

        ip-action-codes
        icmp-action-codes
        tcp-action-codes
        ssl-action-codes
5         http-action-codes )
    )

```

Now, the dispositions and policy rules built into the Policy Engine. These rules can be overwritten by user-defined policy rules.

```

10     ( disposition ok
        ( code OK )
    )

15     ( disposition continue
        ( code CONTINUE )
    )

    ( disposition policy-error
20     ( description "Policy error caused by uncaught event" )
        ( code POLICY_ERROR )
        ( log-directive
            CRITICAL
            "Uncaught event" )
25     )

    ( rule default-rule
        ( description "Catch-all rule for all protocols" )
30     ( protocol present )
        ( action present )
        ( initiator ignore )
        ( target ignore )
        ( outcome
35     ( final
            ( default policy-error )
        )
    )
)
40

```

It is noted that the list of built-in objects included in Table A is by no means complete. In other embodiments, the set of built-in objects is expanded or reduced to reflect the set of protocols supported by the Policy Monitoring System.

5

It is noted that in the preferred embodiment the Policy Engine 101 ranks `default-rule` lower than any user-defined rule. For example, a user-defined rule having initiator and target credentials set to ignore ranks higher than using `default-rule`.

10

In a preferred embodiment, security policy decisions are also affected by any previous history of security violations involving one or both of the principals. Fig. 4 is a schematic diagram of communicating parties according to the invention; wherein an initiator host machine 141 attempts to contact a target host machine 142 over a network, and its events are listened to by an Agent 102 and events are passed onto the invention herein 100. For example, a host machine 141 that repeatedly attempts to perform an illegal operation within a given time window may be blacklisted and rendered incapable of conducting further communication activity within the security domain. In one embodiment, a policy manager maintains a count of security policy violations perpetrated by or against specific principals and provides that information to the Policy Engine 101 as input to a policy evaluation procedure.

Specification Language

25 A security policy is formulated using the Policy manager module's policy specification language (Fig. 1) 108. A preferred embodiment chooses a simplified form of S-expressions as the policy specification language 108. S-expressions are LISP-like parenthesized expressions. The preferred embodiment uses a variant of S-expressions devised by Ron Rivest in R. Rivest, *code and description of S-expressions*, <http://theory.lcs.mit.edu/~rivest/sexp.html>, and used in SPKI/SDSI in C. Ellison, *SPKI Certificate Documentation*, <http://www.clark.net/pub/cme/html/spki.html>. In the preferred embodiment the use of Rivest's S-expressions are restricted in a way such that no empty lists

- are allowed and such that each list must have a type (a byte string) as its first element, denoting the type of the object represented by the list. The use of Rivest's S-expressions is further restricted by only requiring support for the canonical and advanced representations of S-expressions, a preferred embodiment of which is depicted in Table B herein below.

- An advantage of using the canonical representation of S-expressions in the preferred embodiment is for digital signature purposes as well as for relatively efficient communication. It is easy to parse, fairly compact, and is unique for any given S-expression. An advantage of using the advanced representation of S-expressions is for human consumption. It can be thought of as a pretty print of the canonical representation.

- Table B

An example of an advanced representation:

```
( certificate ( issuer alice ) ( subject bob ) )
```

- An example of a canonical representation:

```
(11:certificate(6:issuer5:alice)(7:subject3:bob))
```

- It should be noted that replacing language tokens (e.g. certificate, issuer) with minimally encoded identifiers further optimizes the canonical representation.

The main advantages of using S-expressions in the preferred embodiment are:

- It is easy to represent arbitrary data with S-expressions.
- S-expressions are easy to extend and modify.
- In their advanced representation they are easy to read and edit using, for example, a simple text editor.
- Their canonical representation was designed for efficient packing and parsing. In particular, parsing requires minimal look-ahead and no re-scanning.

- Their canonical representation allows for easy transportation, for example, in files or email messages.
- In their canonical encoding they can be digitally signed.
- It is relatively simple to convert between the advanced and the canonical representation of S-expressions.

A formal description of the policy specification language 108 is provided herein below in Table C.

10

Table C

This table contains a Backus-Naur Form (BNF) description of the grammar for the policy specification language, including an annotation section. All valid policies derive from the <policy> production. This grammar applies to the policy specification after all comments are removed and all macros are expanded. Comments begin with *!//* and extend to the end-of-line. Macros are defined using the C macro syntax.

Incomplete parts of the grammar are noted in *italics*. Terminals are shown in **bold**.

20

```
//
// Basic language stuff
//
```

```
// Terminals requiring further syntax specification
```

25

```
<integer> ::= TBD           // [0-9]*
<symbol> ::= TBD           // alphanumeric and '-',
                           // '_', starts with
```

```
letter
```

30

```
<string> ::= <concat> | TBD // any ASCII character
                           // enclosed in double-
```

```
quotes
```

```
<mac-addr> ::= TBD         // 6 hex byte values
                           // separated by '-'
```

```
<ip-addr> ::= TBD         // IPv4 dotted decimal
                           // notation
```

35

```
<ip-mask> ::= TBD         // address prefix per
                           // RFC-2280
```

```
<hex-string> ::= TBD      // n hex byte values
                           // separated by ':'
```

```

<version-string> ::= TBD           // a string of the form
                                   // <major>.<minor>

// Some productions used only for clarity
5      <name> ::= <symbol>
      <type> ::= <symbol>
      <attr-name> ::= <symbol>

10     // The basic types in the language
      <atom> ::= <symbol> | <integer> | <string> |
               <ip-addr> | <mac-addr> |
               <version> | <hash-atom> | <bool>

15     //
      // Productions for the policy specification section
      //

      // These are values that describe the values of things with values
20     <meta-value> ::= present | absent | ignore

      // Productions used in a few places
      <assertion> ::= ( assertion <bool-expr> )

25     // Version conversion function
      <version> ::= ( version <version-string> )

      // Attributes, used as arguments in predicates and other operations
      <attr-part-list> ::= <atom> |
30      <attr-part-list> <atom>
      <attr-op> ::= ( <attr-name> <attr-part-list> )
      <attribute> ::= <attr-name> | <attr-op>

      // Hashes
35     <hash-alg-name> ::= [Some set of terminals of type hash_alg_t]
      <hash-op> ::= ( hash <hash-alg-name> <attribute> )
      <hash-atom> ::= <hex-string>
      <hash> ::= <hash-atom> | <hash-op>

40     // Operations that operate on attributes and return a
      // basic type (atom)

```

```

    <atom-ops> ::= <hash-op>

// Predicates - used in building boolean expressions
    <generic-compare-op> ::= eq
5      <num-compare-op> ::= gt | lt | ge | le
      <rng-compare-op> ::= range
      <string-compare-op> ::= prefix | substring
      <member-compare-op> ::= member
      <ipmask-compare-op> ::= ip-mask
10     <iprng-compare-op> ::= ip-range
      <cred-match-op> ::= root | has
      <presence-op> ::= present | absent

// Generic argument
15     <arg> ::= <attribute> | <atom> | <atom-ops>

    <generic-cmp-arg> ::= <arg>
    <generic-compare> ::= ( <generic-compare-op>
                           <generic-cmp-arg> <generic-cmp-arg> )
20
    <num-cmp-arg> ::= <attribute> | <integer> | <version>
    <num-compare> ::= ( <num-compare-op>
                       <num-cmp-arg> <num-cmp-arg> ) |
                       ( <rng-compare-op>
                         <num-cmp-arg>
25                         <num-cmp-arg> <num-cmp-arg> )

    <string-cmp-arg> ::= <attribute> | <string>
    <string-compare> ::= ( <string-compare-op>
30                         <string-cmp-arg> <string-cmp-arg> )

    <member-cmp-arg> ::= <arg>
    <member-part-list> ::= <member-cmp-arg> <union> |
                          <member-cmp-arg> <member-cmp-arg>
35     <member-compare> ::= ( <member-compare-op> <member-part-list> )

    <ipmask-compare> ::= ( <ipmask-compare-op>
                          <attribute> <ip-mask> )
    <iprng-cmp-arg> ::= <attribute> | <ip-addr>
40     <iprng-compare> ::= ( <iprng-compare-op>
                          <iprng-cmp-arg>

```

```

                                <iprng-cmp-arg> <iprng-cmp-arg> )

    <cred-match> ::= ( <cred-match-op> <attribute> <cred-name>
5      )

    <presence> ::= ( <presence-op> <attribute> )

    <predicate> ::= <generic-compare> | ' <num-compare> |
                                <string-compare> | <member-compare> |
10      <ipmask-compare> | <iprng-compare> |
                                <cred-match> | <presence>

// Boolean expressions
15      <bool-list-op> ::= and | or
    <bool-monadic-op> ::= not
    <bool> ::= TRUE | FALSE

    <bool-list> ::= <bool-expr> | <bool-list> <bool-expr>
20      <bool-expr> ::= <name> | <bool> | <predicate> |
                                ( <bool-monadic-op> <bool-expr> ) |
                                ( <bool-list-op> <bool-list> )

25      // Descriptions are comments that are carried along with the
    constructs

    <string-list> ::= <string> | <string-list> <string>
    <description> ::= ( description <string-list> )

30      // When we need to break up a string across lines, we use concat to
    // put it back together (just like descriptions)
    <concat> ::= ( concat <string-list> )

35      // Unions are unnamed collections of one or more symbols
    // (with matching types)
    <union-symbol> ::= <atom> | <group-name>
    <union-list> ::= <union-symbol>
40      <union-list> <union-symbol> |
                                ( union <union-list> ) |

```

```

        <union-list> ( union <union-list> )
        <union> ::= ( union <union-list> ) | <union-symbol>

5 // Groups are named and typed unions (all symbols must have one type)
    <group-part-list> ::= <union> | <description> <union> |
        <union> <description>
    <group-name> ::= <name>
    <group> ::= ( group <group-name> <type>
10        <group-part-list> )

// Credentials
    <cred-part> ::= <description> | <assertion>
15    <cred-part-list> ::= <cred-part> | <cred-part-list> <cred-part>
    <cred-name> ::= <name>
    <credential> ::= ( credential <cred-name>
        <cred-part-list> )

20 // Dispositions
    <agent-directives> ::= [Some set of terminals of type
        agent_directive_t]
    <agent-directive-list> ::= <agent-directives> |
25    <agent-directive-list> <agent-directives>
    <agent-directive> ::= ( agent-directive <agent-directive-list> )

    // preliminary list of severity codes; more TBD
    <log-severity> ::= [Some set of terminals of type severity_t]
30    <log-directive> ::= ( log-directive <log-severity> <string> )

        ( log-directive <log-severity> )

    // preliminary list of disposition codes; more TBD
35    <disposition-code> ::= [Some set of terminals of type code_t]
    <disp-code> ::= ( code <disposition-code> )

    <disp-part> ::= <description> | <disp-code> |
        <log-directive> | <agent-directive>
40    <disp-part-list> ::= <disp-part> | <disp-part-list> <disp-part>
    <disposition-name> ::= <name>

```



```

        <disposition> ::= ( disposition <disposition-name>
                               <disp-part-list> )

// Conditions
5   <condition-part-list> ::= <description> <assertion> |
                               <assertion> <description> |
                               <assertion>

        <condition-name> ::= <name>
        <condition> ::= ( condition <condition-name>
10   <condition-part-list> )

// Outcomes are bindings of conditions to dispositions used in rules
        <outcome-default> ::= ( default <disposition-name> )
15   <guard> ::= if | ifnot
        <outcome-clause> ::= ( <guard> <condition-name>
                               <disposition-name> )
        <outcome-list> ::= <outcome-default> |
                               <outcome-clause> <outcome-list>
20   <immediate> ::= ( immediate <outcome-list> )
        <final> ::= ( final <outcome-list> )
        <outcome-types> ::= <immediate> | <final> |
                               <immediate> <final> |
                               <final> <immediate>
25   <outcome> ::= ( outcome <outcome-types> )

// Rules
        <protocol> ::= ( protocol <union> ) |
30   ( protocol <meta-value> )
        <action> ::= ( action <union> ) |
                     ( action <meta-value> )

        <initiator> ::= ( initiator <cred-name> ) |
35   ( initiator <meta-value> )
        <target> ::= ( target <cred-name> ) |
                     ( target <meta-value> )
        <agent> ::= ( agent <cred-name> )
        <rule-name> ::= <name>
40   <rule-list> ::= <rule-name> | <rule-list> <rule-name>
        <prerequisite> ::= ( prerequisite <rule-list> )

```

```

<rank-above> ::= ( rank-above <rule-name> )

<rule-part> ::= <description> | <protocol> | <action> |
               <initiator> | <target> | <outcome> |
5               <prerequisite> | <rank-above> | <agent>
<rule-part-list> ::= <rule-part> | <rule-part-list> <rule-part>
               <rule> ::= ( rule <rule-name>
                           <rule-part-list> )

10
// Policy is the top-level construct of the policy
// specification section
<policy-part> ::= <description> | <group> | <credential> |
                 <condition> | <disposition> | <rule>
15 <policy-part-list> ::= <policy-part> |
                       <policy-part-list> <policy-part>
<policy-name> ::= <name>
<language-version> ::= <version-string>
<policy> ::= ( policy <policy-name> <language-version>
20               <policy-part-list> )

//
// Productions for the annotation section
25 //

<assertion-type> ::= <meta-value> | single-value | multi-value
<weight> ::= ( weight <symbol> <assertion-type> )
<weight-penalty> ::= ( weight-penalty <integer> )
30 <rank-cred-part> ::= <weight> | <weight-penalty>
<rank-cred-part-list> ::= <rank-cred-part> |
                        <rank-cred-part-list> <rank-cred-part>
<ranked-credential> ::= ( ranked-credential <cred-name>
                        <rank-cred-part-list>
35 // Annotation is the top-level construct of the annotation section
<annotation-part-list> ::= <ranked-credential> |
                        <annotation-part-list> <ranked-credential>
<annotation> ::= ( annotation <policy-name>
                    <language-version>
40                    <annotation-part-list> )

```

- In the preferred embodiment the policy specification language 108 is typed. The policy compiler 101 performs the necessary type checking of all S-expressions found in a policy specification 107. Typing aids in catching and diagnosing both common and subtle user errors. A preferred embodiment of the type information is described herein below in Table D.

Table D

Types

- The subtables in this section describe, in a pseudo-formal manner, the type information in the language that is enforced by the parser. The types used should be self-evident.
- First some notation used throughout the tables:

- (list-of type) - (foo (list-of T)) \rightarrow (foo A B C) where A, B, and C are of the type T; the list must have at least one element.
- (multi-of type) - (foo (multi-of T)) \rightarrow (foo (union A B C)) where A, B, and C are of the type T or \rightarrow (foo ABC) where ABC is the name of a group of type T.
- (mix-of type1 type2 type3) - (foo (mix-of R S T)) \rightarrow (foo A B C D) where A, B, C, and D are randomly of types R, S, and T.
- match - a type that is required to be the same as all other occurrences of match in the expression - (foo match match) requires that the two arguments of foo be of the same type (e.g., int_t, string_t).

The following table lists the typed attributes used in conditions and credentials.

Attribute Name	Applicable Protocols	Argument Types	Result Type
agent-attribute	all-protocols	—	(multi-of agent_attr_t)
auth-status	SSL	—	auth_status_t
cert-status	SSL	—	cert_status_t
der-cert	SSL	—	octet_string_t
encipher-keysize	SSL	—	int_t
http-cookie	HTTP	string_t	(multi-of string_t)
http-password	HTTP	—	string_t
http-req-hdr	HTTP	string_t	string_t
http-resp-hdr	HTTP	string_t	string_t
http-set-cookie	HTTP	string_t	(multi-of string_t)

Attribute Name	Applicable Protocols	Argument Types	Result Type
http-status-code	HTTP	—	int_t
http-username	HTTP	—	string_t
icmp-gateway-address	ICMP	—	ip_addr_t
icmp-nested-address	ICMP	—	ip_addr_t
icmp-nested-port	ICMP	—	int_t
initiator-access-rate	all-protocols	—	int_t
initiator-auth-keysize	SSL	—	int_t
initiator-violation-rate	all-protocols	—	int_t
ip-address	IP UDP TCP ICMP	—	ip_addr_t
ip-port	IP UDP TCP ICMP	—	int_t
ke-keysize	SSL	—	int_t
mac-address	IP UDP TCP ICMP	—	mac_addr_t
protocol-version	all-protocols	—	version_t
ssl-ciphersuite	SSL	—	ciphersuite_t
target-access-rate	all-protocols	—	int_t
target-auth-keysize	SSL	—	int_t
target-violation-rate	all-protocols	—	int_t
url	HTTP	—	string_t
x509-cert-path	SSL	—	cert_path_t
x509-issuer	SSL	—	string_t
x509-subject	SSL	—	string_t

The table below lists all the operations in the language that return a dynamic result. For each operation it shows both argument and result types

5

Operation	Result Type	Argument Types
absent	bool_t	string_t
and	bool_t	(list-of bool_t)
default	disposition_t	disposition_t
eq	bool_t	match match
has	bool_t	cert_path_t cred_t
hash	base64_t	hash_alg_t octet_string_t
ge ¹	bool_t	match match
gt ¹	bool_t	match match

¹ Operator only supports types int_t and version_t as arguments.

if	disposition_t	condition_t disposition_t
ifnot	disposition_t	condition_t disposition_t
ip-mask	bool_t	ip_addr_t ip_mask_t
ip-range	bool_t	ip_addr_t ip_addr_t ip_addr_t
le'	bool_t	match match
lt'	bool_t	match match
member	bool_t	match (multi-of match)
not	bool_t	bool_t
or	bool_t	(list-of bool_t)
prefix	bool_t	string_t string_t
present	bool_t	string_t
range'	bool_t	match match match
root	bool_t	cert_path_t cred_t
substring	bool_t	string_t string_t
version	version_t	string_t

The table below is pushing the concept of □type□ far beyond its normal meaning since, in it, we often use type merely to convey positional information. It shows the type of every object in the language and the types of their arguments.

5

Object Name	Object Type	□Argument□ Types
action	act_t	(multi-of action_t)
agent	agt_t	credential_t
agent-directive	agtdir_t	(multi-of agent_directive_t)
assertion	assert_t	bool_t
code	code_def_t	code_t
condition	cond_t	condition_t (mix-of desc_t bool_t)
credential	cred_t	credential_t (mix-of assert_t desc_t prot_t)
description	desc_t	(list-of string_t)
disposition	disp_t	disposition_t (mix-of desc_t code_def_t log_t agtdir_t)
final	dispo_t	(list-of guard_t)
group	group_t	match type_t (multi-of match)
immediate	dispo_t	(list-of guard_t)
initiator	init_t	credential_t
log-directive	log_t	severity_t string_t
outcome	out_t	(list-of dispo_t)
policy	policy_def_t	policy_t string_t (mix-of desc_t group_t cred_t cond_t disp_t rule_def_t)
prerequisite	pre_t	(list-of rule_t)
protocol	prot_t	(multi-of protocol_t)

Object Name	Object Type	Argument Types
rank-above	rank_t	rule_t
rule	rule_def_t	rule_t (mix-of desc_t agt_t prot_t act_t init_t targ_t out_t pre_t rank_t)
target	targ_t	credential_t
union	(multi-of match)	(list-of match)

It is noted that the list of credential and condition attributes included in Table D is by no means complete. In other embodiments, the set of attributes is expanded or reduced to reflect the set of protocols supported by the Policy

5 Monitoring System.

It is noted that although the remainder of this disclosure describes the specification language 108 by means of examples, and that for improved readability, said examples use the advanced rather than the canonical representation of S-expressions, this is not meant to further limit the invention.

10

In the preferred embodiment of the invention, the language 108 allows for comments to be embedded in S-expressions. A comment is allowed anywhere whitespace is valid. A comment begins with "/" and continues to the end-of-line. In compilation, comments are ignored because they serve merely as an aid to the human user.

15

In the preferred embodiment of the invention, the language 108 allows for external files to be included using the #include syntax of C. Included files are supported to enhance modularity and reusability of policy language segments.

20

In the preferred embodiment of the invention, the language 108 allows for macros to be defined using the #define syntax of C. Macros are supported to enhance readability. By convention, macros start with an uppercase letter but need not be fully capitalized.

25

The language 108 comprises the following first-class objects:

- Condition
- Credential
- Disposition

30

- Group
- Policy
- Rule

- 5 In the preferred embodiment first-class objects have names. Names are normally used to refer to an object from another object. By convention, names of built-in objects start with a lowercase letter and use hyphens (-) to separate words. Names of user-defined objects start with an uppercase letter and use intercaps or underscores (_) to separate words, but do not use
- 10 hyphens. Names of data types start with a lowercase letter and end with an underscore followed by a lowercase 't' (_t).

In the preferred embodiment a named object must be defined before its name can be used. The scope of a name is that of the entire policy specification as

15 defined by the policy object.

- In the preferred embodiment first-class objects may optionally include a description field. The description provides human readable text associated with the object. Unlike comments, description fields are preserved by the
- 20 policy parser. When using the advanced representation, description strings may be split across several lines, using the C rules of string concatenation. That is, following the description token are one or more character strings, each enclosed in a set of double quotes.

25 Policy

- In the preferred embodiment a policy is the top-most object defined by the specification language 108 and includes all other first-class objects. A policy manager may load several policies into its internal database. However, at any one point in time, only one active policy is in effect. That is the policy known
- 30 to the Policy Engine 101. Following is an example of a policy object.

```
( policy Sample_Policy_1 "1.0" // policy <name> <version>
  ( description "This is a policy specification description"
    "that is continued on a second line" )
  ...
```

35

)

- In the preferred embodiment a policy object has two mandatory parameters: name, which is used to reference the policy, and version number, which defines the version of the policy specification language 108. A policy's version number is used to check for compatibility between a policy specification and a policy compiler.

Groups and Unions

- 10 In the preferred embodiment groups are named collections of a given type. The union object creates the collection from a set of items. The group object gives the union a name and a type. Following is an example expressing a collection of colors:

```
15  { group SomeColors color_t      // group <name> <type>
    ( description "Some colors I like" )
    ( union RED GREEN YELLOW )
  }
```

- 20 In the example, the object identifies RED, GREEN and YELLOW as items, *i.e.* symbols, of type `color_t` (a fictitious data type) collected in a set named `SomeColors`. By convention, symbols defined in unions are fully capitalized.

- In the preferred embodiment once a symbol is identified as being of a certain type, it is transparently added to an unnamed set of items of that type. It may then be reused in other unions, groups or wherever an individual item of that type is valid. For example, a valid way to define another group is as follows:

```
20  { group RedByAnyOtherName color_t
    ( description "Red in different languages" )
    ( union RED ROSSO ROUGE VERMELHO )
  }
```


However in the preferred embodiment the following group would not be allowed since RED would already have been tagged as being of type color_t.

```

5  ( group AfewOfMyFavoriteThings thing_t
      ( union RED PASTA WINE ) // ERROR! RED previously defined as
    )                          // having type color_t

```

10 In the preferred embodiment sets can be combined with other predefined sets. For example,

```

      ( group MoreColors color_t
        ( union
          SomeColors
15      RedByAnyOtherName      // overlapping ok
          PURPLE BEIGE BURGUNDY
        )
      )

```

20 It is noted that RED overlaps both SomeColors and RedByAnyOtherName, which according to the invention is perfectly acceptable. The resulting set will include only one instance of the set item RED.

25 In the preferred embodiment unions are similar to the C enum type, with the added benefit that unions can be combined and extended without concern for conflicting item values.

30 In a preferred embodiment unions are used, but are not limited to, to define collections of items, such as, for example, IP addresses, MAC addresses, integers, version numbers and hash values. That is, unions can define any data item that has a primitive data type in the language. An example of a group of IP addresses is defined as:

```

35  ( group MyComputers ip_addr_t
      ( union
        207.5.63.23      // desktop at work
        207.5.63.42      // laptop
      )
    )

```

```
128.7.16.64 // home computer
```

```
)
```

```
)
```

- 5 In the preferred embodiment the type of the items in the union must agree with the type specified in the group.

In a preferred embodiment, groups are referenced from other first-class objects. For example, groups are typically used to define collections of
10 protocol actions, SSL ciphersuites, and IP addresses. Note that wherever a group is allowed, the following are also valid:

- A union object (essentially, an unnamed group) provided that any symbols used as elements in the union have already been given a type via a group definition.
- 15 • A single collection item. This is equivalent to a union object with a single element. ~~If the item is a symbol, its type must have been previously defined in a group.~~

A list of built-in groups is given in section Table A.

20

Credentials

In the preferred embodiment a credential is a statement about a principal in a protocol event. It consists of a logical expression containing one or more assertions about the attributes that make up a principal's credentials. When
25 a policy rule is evaluated against a protocol event, the credentials presented in the protocol event are compared to the credentials specified in a purported credential object. If the logical expression defined in the credential object is satisfied, the principal's presented credentials are said to satisfy the purported credentials. As an example, the following purported credentials are satisfied if
30 the principal's IP address is 207.5.63.8 and its IP port number is either 80 or greater than 443.

```
{ credential Credentials_Example_1 // credential <name>
  { assertion
    { and
```

35

```

    ( eq ip-address 207.5.63.8 )
    ( or
      ( eq ip-port 80 )
      ( gt ip-port 443 )
    )
  )
)

```

10 In the preferred embodiment each protocol has a set of attributes that may be used to build purported credentials. Table E herein below lists all the attributes currently defined and, for each attribute, it shows the protocols where the attribute might be included in the presented credentials, as well as the operations where the attribute may be used as an operand.

15

Table E

Attribute Name	Applicable Protocols	Description	Compare Operations
<i>agent-attribute</i>	all-protocols ²	The attributes of the reporting Agent, as a union of symbolic names	member
<i>cert-status</i>	SSL	The validity status of a certificate	eq, member
<i>der-cert</i>	SSL	A DER encoded certificate	hash
<i>http-password</i>	HTTP	The password used in basic authentication	eq, member, substring, prefix
<i>http-username</i>	HTTP	The user name used in basic authentication	eq, member, substring, prefix
<i>ip-address</i>	IP UDP TCP ICMP	An IP address	eq, member, ip-mask, ip-range
<i>ip-port</i>	IP UDP TCP ICMP	An IP port	eq, member, gt, ge, lt, le, range
<i>mac-address</i>	IP UDP TCP ICMP	A MAC address	eq, member
<i>url</i>	HTTP	A URL	eq, member, substring, prefix
<i>x509-cert-path</i>	SSL	An X.509 certificate chain	root, has
<i>x509-issuer</i>	SSL	An X.509 certificate issuer	eq, member, substring,

² Can be used to identify the reporting Agent in any policy rule but must not be mixed with other credential attributes.

Attribute Name	Applicable Protocols	Description	Compare Operations
			prefix
x509-subject	SSL	An X.509 certificate subject	eq, member, substring, prefix

It is noted that the list of credential attributes included in Table E is by no means complete. In other embodiments, the set of attributes is expanded or reduced to reflect the set of protocols supported by the Policy Monitoring System.

In the preferred embodiment each attribute can be thought of as having an implied getter function that returns its value. Most attribute getters take no arguments and return a single value. In the preferred embodiment, however, some attribute getters (e.g. `http-req-hdr` and `http-cookie`) are functions that take one or more arguments and may return complex results. For example, `http-cookie` takes as an argument the name of a cookie in an HTTP request header and returns its value or values as a union of strings.

In the preferred embodiment it is important not to mix credential attributes from different protocol sets in a credential specification. For example, combining `ip-address` and `der-cert` in the same credential object would be an error and flagged by the policy compiler. As another example, using a credential in a policy rule for a protocol action that is incompatible with the credential attributes in the credential object is considered an error, flagged by the policy compiler. However, it is possible to use those attributes in two separate credential objects and establish relationships between them within policy rules (e.g. access to resource X is restricted to principals previously authenticated with credentials Y). See example `Check_Access_Denial` herein below for an example of establishing this type of relationships in policy rules.

In the preferred embodiment the credential attribute `agent-attribute` is used to define the credentials of the Agent 102 reporting the protocol event 103. Agents are individually configured with a set of attributes, which are used to

identify them to a policy manager. In another embodiment, some agent attributes might uniquely identify a specific Agent (e.g. MONITOR_NEXT_TO_ROUTER_X) while others might identify a group of Agents (e.g. ALL_MONITORS_IN_SUBNET_Y).

- 5 The agent-attributes attribute returns a union of identification attributes for the reporting Agent 102. In the preferred embodiment within a credential specification, assertions about agent attributes may not be mixed with assertions about any other credential attributes.
- 10 Table F herein below lists all the operations used in a preferred embodiment to make assertions about attributes.

Table F

Operation	Description
<i>absent</i>	Whether (true) the attribute denoted by the operand does not have a value in the protocol event
<i>and</i>	Logical AND of a list of boolean expressions, its operands
<i>eq</i>	Whether (true) two operands have the same value
<i>ge</i>	Whether (true) the first operand's value is greater than, or equal to, the second's
<i>gt</i>	Whether (true) the first operand's value is greater than the second's
<i>has</i>	Whether (true) the certificate chain defined by the first operand's value contains a certificate that satisfies the second operand (a credential name)
<i>hash</i>	Computes a digest of the second operand's value using the hashing function defined by the first operand; returns the hash as a hexadecimal string
<i>ip-mask</i>	Whether (true) the first operand's value is included in the set of IP addresses defined by the second operand, an IP address prefix [RFC2280]; an address prefix is represented as an IPv4 address (dotted-decimal format with four integers) followed by the character slash '/' followed by an integer in the range from 0 to 32. The latter denotes the number of high-order bits from the preceding address that constitute a subnetwork address. If the subnetwork address bits match exactly the corresponding bits in the first operand's value, the operation returns true. The following are valid address prefixes: 128.9.128.5/32, 128.9.0.0/16, 0.0.0.0/0; the following address prefixes are invalid: 0/0, 128.9/16 since 0 or 128.9 are not dotted-decimal strings containing four integers.

Operation	Description
<i>ip-range</i>	Whether (true) the first operand's value is included in the set of IPv4 addresses defined by an IP address range whose lower bound is the operand with the lower value and whose upper bound is the operand with the higher value; the three operand values are taken as 32-bit unsigned integers and, if the first operand value falls within the inclusive numerical range defined by the two other operand values, the operation returns true
<i>le</i>	Whether (true) the first operand's value is less than, or equal to, the second's
<i>lt</i>	Whether (true) the first operand's value is less than the second's
<i>member</i>	Whether (true) the first operand's value is a member of the set defined by the second operand (a union)
<i>not</i>	Logical negation of its operand's value
<i>or</i>	Logical OR of a list of boolean expressions, its operands
<i>prefix</i>	Whether (true) the string that constitutes the first operand's value includes, starting at the first character, the string defined by the second operand
<i>present</i>	Whether (true) the attribute denoted by the operand has a value in the protocol event
<i>range</i>	Whether (true) the first operand's value is within the inclusive numerical range defined by the values of the second and third operands; the range comprises the set of values between the lower operand value and the higher
<i>root</i>	Whether (true) the certificate chain defined by the first operand's value has, as its root, a certificate that satisfies the second operand (a credential name)
<i>substring</i>	Whether (true) the string that constitutes the first operand's value includes the string defined by the second operand

It is noted that the list of operations included in Table F is by no means complete. In other embodiments, the set of operations is expanded or reduced to reflect the set of protocols and features supported by the Policy Monitoring System.

In the preferred embodiment credentials may be combined with other credentials or with additional assertions. Consider the following example:

```

10  ( credential Credentials_Example_2
      ( assertion
        ( or
          Credentials_Example_1
          ( and
15      ( ip-mask ip-address 207.5.0.0/16 )
        ( range ip-port 25 443 )
        )
      )
    )
  )

```

```

    )
  )
)

```

- 5 The example herein above defines purported credentials that will be satisfied if either `Credentials_Example_1` is satisfied or if the presented credentials' IP address falls within the subnetwork defined by the address prefix 207.5.0.0/16 and if the IP port is between 25 and 443, inclusive.
- 10 In the preferred embodiment the absence of an assertion about a specific attribute in a credential specification indicates that its value is to be ignored in considering the presented credentials. In the preferred embodiment, it is often useful to indicate that a particular attribute must or must not be specified in the presented credentials, irrespective of the attribute's value, if any. The
- 15 operations absent and present accomplish this, as illustrated by the following examples:

```

( credential Credentials_Example_3
  ( assertion
20    ( and
      // http-username must exist, but don't care about its
value
      ( present http-username )
      // the absence of an assertion about http-password
25    indicates
      // that its presence or absence is irrelevant
    )
  )
)
30
( credential Credentials_Example_4
  ( assertion
    ( and
      // an X.509 certificate must not have been presented
35    ( absent der-cert )
    )
  )
)

```

Conditions

In the preferred embodiment a condition defines a constraint upon a protocol event 103. Said condition comprises a logical expression containing one or more assertions about attributes of the protocol event. Policy rules use conditions to specify particular constraints that must or must not be satisfied by the protocol event 103.

Table G lists attributes of a protocol event 103 that may be used when formulating conditions. For each attribute the table shows protocols for which the attribute is defined, as well as the operations which can take the attribute as an operand.

Table G

Attribute Name	Applicable Protocols	Description	Compare Operations
<i>auth-status</i>	SSL	The status of an authenticated session at the end of the authentication handshake	eq, member
<i>encipher-keysize</i>	SSL	The size of the key used for data encipherment (e.g., size of an IDEA key)	eq, member, gt, ge, lt, le, range
<i>http-cookie</i>	HTTP	Takes as an argument the name of a cookie in the request header and returns its value(s) as a union of strings	member
<i>http-req-hdr</i>	HTTP	Takes as an argument the name of a client request header and returns its value	eq, member, substring, prefix
<i>http-resp-hdr</i>	HTTP	Takes as an argument the name of a server response header and returns its value	eq, member, substring, prefix
<i>http-set-cookie</i>	HTTP	Takes as an argument the name of a cookie in the response header and returns its value(s) as a union of strings	member

Attribute Name	Applicable Protocols	Description	Compare Operations
<i>http-status-code</i>	HTTP	The status code returned on HTTP responses (aka response code)	eq, member, gt, ge, lt, le, range
<i>icmp-gateway-address</i>	ICMP	The IP address of the gateway host on a redirect message	eq, member, ip-mask, ip-range
<i>icmp-nested-address</i>	ICMP	The IP address carried in a [destination unreachable] message	eq, member, ip-mask, ip-range
<i>icmp-nested-port</i>	ICMP	The port number carried in a [destination unreachable] message	eq, member, gt, ge, lt, le, range
<i>initiator-access-rate</i>	all-protocols	The rate at which the current active principal has been the initiator of communications, over a predefined (configurable) period of time	eq, member, gt, ge, lt, le, range
<i>initiator-auth-keysize</i>	SSL	The size of the key used for initiator authentication and/or digital signatures (e.g., size of public key modulus)	eq, member, gt, ge, lt, le, range
<i>initiator-violation-rate</i>	all-protocols	The rate at which the current active principal has been the initiator of security policy violations, over a predefined (configurable) period of time	eq, member, gt, ge, lt, le, range
<i>ke-keysize</i>	SSL	The size of the key-encipherment key (e.g., size of public key modulus)	eq, member, gt, ge, lt, le, range
<i>protocol-version</i>	all-protocols	The version of the protocol	eq, member, gt, ge, lt, le, range
<i>ssl-ciphersuite</i>	SSL	The negotiated ciphersuite	eq, member
<i>target-access-rate</i>	all-protocols	The rate at which the current passive principal has been the target of communications, over a predefined (configurable) period of time	eq, member, gt, ge, lt, le, range
<i>target-auth-keysize</i>	SSL	The size of the key used for target authentication and/or digital signatures (e.g., size of public key modulus)	eq, member, gt, ge, lt, le, range
<i>target-violation-rate</i>	all-protocols	The rate at which the current passive principal has been the target of security policy violations, over a predefined (configurable) period of time	eq, member, gt, ge, lt, le, range

It is noted that the list of condition attributes included in Table G is by no means complete. In other embodiments, the set of attributes is expanded or reduced to reflect the set of protocols and features supported by the Policy Monitoring System.

5

In the preferred embodiment operations listed in Table G may be used to build assertions about condition attributes.

In the preferred embodiment condition attributes cannot mix with those from different protocol sets in a condition specification. A condition used in a policy rule for a protocol that is incompatible with the condition attributes in the condition object is considered an error and is flagged by the policy compiler. For example, it is illegal to use `ssl-ciphersuite` in a condition referenced by a policy rule for HTTP.

15

Following are some examples:

```
( group Strong_RSA_Ciphersuites  ciphersuite_t
  ( description "Strong ciphers with RSA key exchange" )
20  ( union  SSL_RSA_WITH_RC4_128_MD5
        SSL_RSA_WITH_RC4_128_SHA
        SSL_RSA_WITH_IDEA_CBC_SHA
        SSL_RSA_WITH_3DES_EDE_CBC_SHA
        SSL_DH_RSA_WITH_3DES_EDE_CBC_SHA
25  SSL_DHE_RSA_WITH_3DES_EDE_CBC_SHA
    )
  )

( condition SslV3StrongCiphers          // condition <name>
30  ( assertion
    ( and
      ( ge protocol-version ( version "3.0" ) )
      ( member ssl-ciphersuite Strong_RSA_Ciphersuites )
      ( ge ke-keysize 768 )
35  ( ge target-auth-keysize 1024 )
    )
  )
)
```

```

    ( condition HackerTripwire
      ( assertion
        ( ge initiator-violation-rate 10 )
5      )
    )

    ( condition ProtectSSL
      ( assertion
10      ( and SslV3StrongCiphers HackerTripwire )
      )
    )

```

- Herein above, the condition `SslV3StrongCiphers` can be used with an
- 15 SSL protocol event to ensure that SSL 3.0 or higher is used, that the negotiated ciphersuite is one of the strong RSA-based ciphersuites, that the RSA key-encipherment key has a modulus of no less than 768 bits, and that the RSA authentication key has a modulus of no less than 1024 bits.
- 20 Herein above, the condition `HackerTripwire` can be used with any protocol event 103 to ensure that the active principal 141 is not a potential attacker. The third condition, `ProtectSSL`, simply combines the first two.

Dispositions

- 25 In the preferred embodiment a disposition defines an outcome of a policy rule. Each policy rule may have many possible outcomes depending on, for example, constraints imposed on the protocol event.

- 30 See Table H herein for a list of disposition codes and an explanation of their meanings in the preferred embodiment.

Table H

Disposition Code	Description
OK	The network event conforms to the security policy

Disposition Code	Description
CONTINUE	Additional information is needed before determining whether or not the network event conforms to the security policy
ACCESS_DENIED	Access to the target resource is denied by the security policy
AUTHENTICATION_VIOLATION	Authentication between the communication parties does not conform to the requirements set out by the security policy
SECURITY_ATTACK	A security attack has been detected
SECURITY_QOS	The security quality of service parameters associated with a protocol event do not meet the requirements set out by the security policy
POLICY_ERROR	An error has been detected in the security policy specification

It is noted that the list of disposition codes included in Table H is by no means complete. In other embodiments, the set of disposition codes is expanded or reduced to reflect the set of features supported by the Policy Monitoring System.

Table I herein below lists possible severity codes in the preferred embodiment.

Table I

Severity Code	Description
CRITICAL	Critical security violation, e.g., the network is undergoing an active security attack
HIGH	High-severity security violation, e.g., attempt to access sensitive data
MEDIUM	Medium-severity security violation, e.g., attempt to access a protected (but not highly sensitive) resource
WARNING	Low-severity security violation, e.g., an incorrect password was entered
MONITOR	A security violation was not detected but an unusual or potentially suspect network event has occurred, e.g., TELNET access to a public web server

Severity Code	Description
INFORMATION	A perfectly valid network event is being reported for informational purposes only

It is noted that the list of severity codes included in Table I is by no means complete. In other embodiments, the set of severity codes is expanded or reduced to reflect the set of features supported by the Policy Monitoring System.

Table J herein below lists possible agent directives in the preferred embodiment.

Table J

Agent Directive	Description
DECRYPT	The Agent is instructed to decrypt all traffic at the current protocol layer
DISRUPT	The Agent is instructed to terminate and/or disrupt all subsequent traffic associated with this network event
LOG_TRAFFIC	The Agent is instructed to log all traffic at the current protocol layer

It is noted that the list of agent directives included in Table J is by no means complete. In other embodiments, the set of agent directives is expanded or reduced to reflect the set of features supported by the Policy Monitoring System.

Following are examples of preferred embodiments of dispositions:

```

20 // Network event ok but should be logged
   ( disposition Ok_Monitor           // disposition <name>
     { code OK }                     // disposition code
     { log-directive                 // logging directive
       MONITOR                      // severity code
25     "Monitored activity" )        // logging string
   )

```

The `Ok_Monitor` disposition is used to dispose of a valid network event 103 while flagging a logging subsystem that this event should be logged at a low severity level (MONITOR).

```

5  // Decrypt SSL session data and continue processing network event
   ( disposition Continue_Decrypt
     ( code CONTINUE )
     ( agent-directive DECRYPT )
   )

```

- 10 The `Continue_Decrypt` disposition is used to inform the Policy Engine 101 that additional information is needed from the Agent 102 before determining a final disposition 105 for the network event 103 while, at the same time, instructing an appropriate Agent to decrypt all traffic at a current protocol layer.

```

15 // access to target resource is denied
   ( disposition Access_Denied
     ( code ACCESS_DENIED )
     ( log-directive
20       HIGH
       "Access denied" )
   )

```

The `Access_Denied` disposition is used as a final disposition 105 for a network event 103. It denotes a policy violation.

- 25 A list of built-in dispositions of the preferred embodiment is provided herein above in Table A.

Rules

- 30 In the preferred embodiment a rule object defines a policy rule. A policy rule governs a specific interaction, or set of interactions, between two communicating entities. The Policy Engine 101 evaluates policy rules against protocol events to determine if the latter conform to the active security policy.
- 35 Following is an example of a policy rule according to a preferred embodiment of the invention:

```

( rule Tcp_Ext2Int                                     // rule <name>
  ( description "Communications from external hosts" )
  ( agent Foo_Subnet_Monitor )                         // the reporting agent
5  ( protocol TCP )                                     // the protocol
  ( action CONNECT )                                   // the protocol action
  ( initiator External_Hosts )                         // the active principal
  ( target Internal_Hosts )                           // the passive principal
  ( outcome
10  ( immediate                                         // the immediate outcome
    // if/ifndef <condition> <disposition>
    ( if Catch_Suspect Security_Attack_Possible )
    ( if Catch_Attacker Security_Attack_Progress )
    ( default continue )                             // using built-in disposition
15  )
    ( final                                           // the final outcome
      ( default Ok_Monitor )
    )
  )
20 )

```

In the preferred embodiment a policy rule comprises:

- 25 • *Agent* – represents Agent 102 that reported the protocol event 103. The Agent 102 is denoted by a credential name. The policy rule is only considered if this credential is satisfied by the credentials presented by the reporting Agent 102. In the example above, *Foo_Subnet_Monitor* is the name of a credential object identifying one or more Agents. This field is optional. If omitted, the rule applies to all Agents.
- 30 • *Protocol* – a protocol to which the rule is applicable. A protocol event 103 addresses one and only one protocol. This field is mandatory. Note that the special token *ignore* is used to denote a rule that applies to all protocols.
- 35 • *Action* – a protocol action to which this rule is applicable. Each protocol comprises one or more several distinct actions (e.g. connect, transfer-data, release), some of which might be of interest to the security policy. A

protocol event denotes one and only one protocol action. This field is mandatory. Note that the special token *ignore* is used to denote a rule that applies to all actions within the specified protocol.

- 5 • *Initiator* – represents the active principal 141 in the protocol event 103. The initiator 141 is denoted by a credential name or by the special tokens *absent* (credentials must not be presented in the protocol event), *present* (credentials must be presented but their actual value is unimportant) and *ignore* (credentials may or may not be presented). In
10 the example herein above, *External_Hosts* is the name of a credential object identifying one or more TCP/IP hosts. This field is mandatory.
- *Target* – represents the passive principal 142 in the protocol event 103. The target 142 is denoted by a credential name or by the special tokens
15 *absent*, *present* and *ignore*. In the example above, *Internal_Hosts* is the name of a credential object identifying one or more TCP/IP hosts. This field is mandatory.
- *Prerequisite* – (not shown in the example above) one or more rules that
20 must be satisfied by a previous protocol event. Prerequisite rules are identified by names. Prerequisites are used to place additional constraints on an entire network event 103. See an example herein that illustrates the use of prerequisites in rules. It should be noted that if two or more rules are listed as prerequisites, the prerequisite is satisfied if *any* of the listed
25 rules taken in the order in which they are listed satisfies a previous protocol event. This field is optional.
- *Outcome* – the outcome section defines what to do with the protocol (or
30 network) 103 event if the current policy rule is applied to the protocol event. That is, if the rule is selected by the Policy Engine 101 as the most suitable for the protocol (or network) event. Every policy rule must have a disposition that applies to the protocol event and another disposition that applies to the entire network event. In some cases these are one and the same. The Policy Engine 101 evaluates the outcome and produces a

disposition for either the protocol or the network event. There are two outcomes defined:

- 5 • *Immediate* – an immediate outcome applies to the protocol event immediately. A policy rule may or may not include an immediate outcome. If it does, the outcome is evaluated as soon as the rule is selected for the protocol event. If it does not, there is an implied disposition for the protocol event, a built-in disposition *continue* (see Table A for the definition) which instructs the Policy Engine 101 to
10 continue processing the network event. If the immediate outcome generates a disposition with a disposition code other than *CONTINUE*, this disposition becomes the disposition for the entire network event. In this instance, the final outcome, defined herein below, will not be evaluated.
- 15 • *Final* – an outcome that applies to the entire network event if this rule becomes a final rule evaluated for that event. The final outcome must be specified if the immediate outcome does not generate a final disposition for the network event. If it is not, an implied disposition for
20 the network event, the built-in disposition *policy-error*, see Table A for the definition, denotes a policy specification error. The final outcome is evaluated when the Policy Engine determines that no additional protocol events are to be considered for the current network event. The final outcome must always generate a final disposition, *i.e.*
25 a disposition with a disposition code of *CONTINUE* is not allowed in a final outcome.

In the preferred embodiment each outcome section comprises one or more conditional statements, each followed by a disposition. The purpose
30 of conditional statements is to specify constraints upon a protocol event, or special conditions that, if satisfied, cause the generation of an alternate disposition for the protocol (or network) event. Conditional statements are evaluated in the order in which they are specified within the outcome section.

In the preferred embodiment a conditional statement starts with one of the following keywords:

- 5 • **if** – takes as arguments a condition and a disposition, each referenced by name. If the condition evaluates to TRUE, the disposition becomes the disposition for the protocol event.
 - **ifnot** – takes as arguments a condition and a disposition, each referenced by name. If the condition evaluates to FALSE, the disposition becomes the disposition for the protocol event.
 - 10 • **default** – takes a single argument, a disposition referenced by name. It is equivalent to a condition that is always satisfied, thereby triggering the disposition that is its argument. This conditional statement is mandatory and must be the last conditional statement in an outcome.
- 15 The following examples illustrate the use of prerequisites in rules in a preferred embodiment. The first rule is the prerequisite.

```

( credential Host_A
  ( assertion
    ( and
      ( eq ip-address 207.5.63.8 )
      ( eq ip-port 80 )
    )
  )
)

( rule Access_Host_A
  ( protocol TCP )
  ( action CONNECT )
  ( initiator ignore )
  ( target Host_A )
  ( outcome
    ( final
      ( default Access_Denied ) // Access_Denied defined above
    )
  )
)

```

Herein above, the rule `Access_Host_A` states that access to host A on port 80 by any host is denied, unless explicitly allowed by a rule at a higher protocol layer. Note the use of a final outcome, which is only evaluated if `Access_Host_A` becomes the applicable rule for the entire network event.

- 5 The implied disposition for the protocol event is CONTINUE.

This rule can be overridden by another rule at the HTTP layer stating that access is allowed to host A on port 80, as shown below:

```

10 ( rule Http_To_Host_A
    ( protocol HTTP )
    ( action ignore )
    ( initiator ignore )
    ( target ignore )
15 ( prerequisite Access_Host_A ) // reference to rule above
    ( outcome
      ( immediate
        ( default ok ) // using built-in disposition
      )
20 )
  )

```

The end result of the two policy rules herein above is to prevent all access to host A on port 80 unless that access is using HTTP over TCP/IP.

- 25 In the preferred embodiment a prerequisite rule is any rule that is selected for a previous protocol event. This includes rules in the same protocol layer. As an example, to ensure that a web server requires HTTP authentication before allowing access to a specific web page, use the following rules:

```

30 ( credential Some_Url
    ( assertion
      ( prefix url "//myserver.com/Documents" )
    )
35 )

( rule Host_A_Anon_Access

```

```

        ( protocol HTTP )
        ( action ( union GET POST ) )
        ( initiator absent )
        ( target Some_Url )
5      ( prerequisite Access_Host_A )    // from example above
        ( outcome
          ( final
            ( default Access_Denied ) // Access_Denied defined above
          )
10     )
    )

    ( condition Require_Auth
      ( description "Check if server returned the Unauthorized response
15         " "code" )
      ( assertion
        ( eq http-status-code 401 )
      )
    )
20

    ( rule Check_Access_Denial
      ( protocol HTTP )
      ( action RESPONSE )
      ( initiator ignore )
25     ( target ignore )
      ( prerequisite Host_A_Anon_Access )
      ( outcome
        ( immediate
          ( ifnot Require_Auth Access_Denied )
30         ( default ok ) // using built-in disposition
        )
      )
    )
  )

```

35 The example herein above shows that access to the document sub-tree identified by Some_Url requires the user be authenticated using basic HTTP authentication. The authentication is accomplished by means of the condition Require_Auth which, in the context of rule Check_Access_Denial, checks that the server returns an *Unauthorized* status code. If the server fails

40 to do so, the Access_Denied disposition is generated. Note that the

prerequisite constraint ensures that the rule `Check_Access_Denial` is only considered if the rule `Host_A_Anon_Access` is selected when the HTTP request event is evaluated, that is, requests where basic HTTP authentication is not used.

5

The Policy Specification Process

In the preferred embodiment the policy specification process comprises the following steps:

- 10 1) Identify communicating entities recognized by the security policy. The entities comprise physical networks and sub-networks, host machines, communication protocols, users, applications, services, and any other resources of interest.
- 15 2) Identify relationships between the communicating entities and define rules to control said relationships (e.g. host A may communicate with host B but not with host C).
- 20 3) Formally define communicating entities and entity relationships using the policy specification language (Fig. 1; 108) according to the invention. In a preferred embodiment a visual tool is used. In another embodiment a text-based editor is used. In the preferred embodiment the output of this step is a policy specification in an advanced encoding format according to the invention.
- 25 4) Compile the policy specification with a Policy Compiler (Fig. 1, 106). In one embodiment, said compilation step is incorporated into a graphical policy editor, such that it is incorporated into said policy specification step. In another embodiment it is a distinct step. This step comprises:
 - 30 a) Checking the specification for errors of syntax or semantics;
 - b) Checking the specification of credentials for errors (e.g. credentials that can never be satisfied);
 - c) Checking the specification of conditions for errors (e.g. conditions that can never be satisfied);

- d) Checking the specification of rules for completeness and coverage;
- e) Ordering credentials based on their specificity (described in detail herein below);
- f) Ordering rules based on the credentials of their principals (described in detail herein below); and
- g) Resulting in an annotated policy specification (Fig. 1, 109) represented by a text file (Fig. 1 107).

The annotated policy specification 107 is suitable for loading into the Policy Engine 101 for evaluation of one or many network events 103, or back into the graphical policy editor for visualization and further refinement.

Evaluation of Rules

This section describes how policy rules are organized and evaluated according to the invention.

15 Policy Evaluation Model

The policy specification language 108 alone does not describe how the Policy Engine 101 evaluates policy rules. In the preferred embodiment of the invention, a security administrator that writes the policy specification 107 and the Policy Engine 101 that enforces the policy specification 107 share a common view of the evaluation procedure. The evaluation of policy rules is deterministic.

In the preferred embodiment of the invention the basic policy specification language 108 is augmented to convey information about how rules are ordered for purposes of evaluation, i.e. which rules are evaluated first and which rules are selected for any given network event. The augmented language is a superset of the basic specification language 108 and it is hereinafter referred to as the annotated specification language 109.

In one embodiment the security administrator uses the annotated specification language 109 using a visual tool, such as a graphical policy editor to determine how the policy rules are interrelated, their hierarchical relationships and how they will be evaluated. This step is crucial to determining whether

the specified policy correctly reflects the desired security policy and to identifying areas where the policy specification needs refinement.

5 In the preferred embodiment the Policy Engine 101 uses the annotated language 109 to organize the policy, after having converted it to an internal representation in a manner best suited for the efficient evaluation of network events.

10 In the preferred embodiment the Policy Engine 101 receives protocol events in proper sequence. Protocol events for protocols lower in the protocol stack are received before protocol events for protocols higher in the stack. This sequencing is important because the Policy Engine 101 must make a policy decision about, for example, a TCP connection, before it makes a decision about an SSL session that uses that TCP connection.

15 Data about a specific protocol event may not arrive all at once. For example, when evaluating an SSL session the Policy Engine 101 first receives the server certificate and negotiated ciphersuite before receiving a client certificate or a message indicating that none was provided. In a preferred
20 embodiment, the Policy Engine 101 uses incomplete information about a protocol event in order to collect a set of possible policy rules applicable to that event. However, for the sake of simplicity, the remainder of this document assumes that Agents convey information about protocol events in an atomic manner.

25 In the preferred embodiment for every protocol event the Policy Engine 101 selects a policy rule applicable to that event. Every policy rule is associated with a specific protocol and action or a set of protocols and actions. Therefore only the set of rules relevant to the protocol event is considered. Of that set,
30 several rules can be satisfied by the event. In the preferred embodiment a policy rule is satisfied by a protocol event if the following holds true:

- 1) The credentials of the Agent 102 reporting the event match the rule's agent credentials (if any), defined as a set of attribute-value assertions.

- 2) The rule's protocol specifier matches the protocol identifier in the protocol event.
 - 3) The rule's action specifier matches the action identifier in the protocol event.
 - 5 4) The rule's prerequisite clause is satisfied (details described herein below).
 - 5) The credentials of the initiator 141 and target 142 principals in the protocol event satisfy the rule's corresponding credentials, defined as a set of attribute-value assertions.
- 10 In the preferred embodiment when several rules are satisfied by a protocol event, the Policy Engine 101 selects a rule that is most specific to the protocol event. The specificity of a policy rule is determined by the specificity of the credentials associated with the policy rule, as well as the specificity of the rule's protocol, action and prerequisite specifiers. For example, a rule that
- 15 targets one single protocol is more specific than a rule that targets all protocols. In another example, a rule that specifies a prerequisite is more specific than a rule that does not.

20 In the preferred embodiment the specificity of a credential specification is determined by the set relationships of said specification with other credential specifications. Following are examples of credential specifications:

- A: all principals with blue eyes
- B: all principals with blue eyes and black hair
- 25 • C: all principals with black hair

B defines the intersection of A and C, i.e. B is a subset of both A and C. Thus, B is more specific than either A or C.

- 30 According to the invention, in general, the more data described about a principal the more specific are the credentials. In the preferred embodiment, some attributes of a principal's credentials have more importance than do other attributes of the credentials. In the preferred embodiment the importance of an attribute is represented by its weight. The attribute weight is
- 35 determined by its role as a discriminator of principals. For example, an

attribute that yields a small set of principals has more weight than an attribute that yields a larger set of principals. In the hair and eye color example herein above, it is arbitrary to give a higher weight to eye color versus hair color or to give both hair and eye color the same weight. Assigning an attribute weight is easier because typically protocol credentials are structured hierarchically. For example, in the TCP protocol, the IP address attribute has clearly more weight than the IP port attribute because the number of principals with a given IP address is generally much smaller than the set of principals with a given port number.

In the preferred embodiment attributes that comprise a set of credentials are ranked by weight and the combined weight of all attributes in a credential specification is considered in determining a relative specificity of said specification.

In the preferred embodiment a policy specification has sets of credentials each of which are ranked at a same specificity level, thereby rendering many policy rules that are applicable to a given protocol event. Herein below is provided a section describing a number of practical guidelines for good policy development that minimize herein above ambiguities.

Fig. 5a is a schematic diagram of the preferred embodiment in which a network event 103 comprises M protocol events at different protocol layers, and in which the network event 103 has an associated network event disposition 105.

Fig. 5b is an algorithm showing how the M protocol events at different protocol layers of the network event 103 result in pending rules with or without immediate outcomes and, finally, a final disposition for the network event 105.

For clarity, the algorithm assumes that the Policy Engine 101 always finds a policy rule applicable to a given protocol event, that at least a first Protocol Event (1) exists, and that the algorithm ends when the Agent 102 informs the Policy Engine 101 that no further protocol events will be generated. These assumptions are for clarifying purposes only and do not limit the invention in any way.

The algorithm begins with $j=1$ (500) and with the Policy Engine 101 receiving Protocol Event (1) from the Agent 102 (501) (502).

- 5 Once a most specific policy rule is selected for a given protocol event (503), the Policy Engine 101 consults an outcome clause (504) determining if an immediate outcome is applied to the protocol event. In the preferred embodiment an immediate outcome applies to a protocol event while a final outcome applies to a network event (103).

10

In the preferred embodiment an immediate outcome is executed when it is specified. The immediate outcome can evaluate constraints (*i.e.* conditions) against a protocol event, produce a set of agent directives (*e.g.* instructing the Agent 102 to decrypt all subsequent traffic), and produce a final disposition (506) for the protocol event rendering said disposition for the entire network event. When a disposition of an immediate outcome is not a final disposition, a special disposition code, CONTINUE, is used as an indicator. All disposition codes other than CONTINUE denote final dispositions.

15

- 20 In the preferred embodiment when an immediate outcome does not produce a final disposition the associated selected policy rule becomes a pending policy rule for the related network event (507). The Policy Engine 101 then waits for further protocol events of the network event 103 from the Agent 102 (508) and (501). In this embodiment, said pending policy rule is overridden by
25 subsequent policy rule selected for a protocol event higher in the associated protocol stack (507).

- In the preferred embodiment policy evaluation ends in one of two cases. First case is when no further rules in the policy apply to a network event (*e.g.* a highest protocol in the stack is reached). Second case is when the Agent 102
30 informs the Policy Engine 101 that no further protocol events will be generated (502) (509) (506) (510). In either case, a policy decision is then expected for the entire network event. The Policy Engine 101 selects a pending policy rule for a protocol highest in the protocol stack and executes
35 the final outcome defined for that rule (511). In the preferred embodiment

constraints are evaluated against the entire network event. In the preferred embodiment a final outcome always produces a final disposition (509) which becomes a disposition for the network event (506).

- 5 In the preferred embodiment a protocol event must result in a selection of a policy rule (pending or final). When a policy rule applicable to a given protocol event is not found, the Policy Engine 101 produces a special disposition identifying a policy specification error. See the default policy rule in Table A.

10 *Ordering of Credentials*

In the preferred embodiment credentials are ordered based on a combined weight of all attribute-value assertions that make up a credential specification.

- 15 In the preferred embodiment computing a weight of an attribute-value assertion of an attribute requires the following two steps:

- 20 1) Assigning a ranking value to the attribute. Attributes that are listed in a credential specification are ranked against each other. Ranking is based on a value of the attribute as a discriminator of principals identified by the credentials. If the presence of the attribute in a credential specification generally yields a smaller set of principals than the presence of another attribute, then the former has a higher ranking than the latter.

- 25 2) Assigning a ranking value to an assertion type of the attribute. An assertion type is used to make an assertion about the value of an attribute (e.g. eq, substring, range). Following are five assertion types, in decreasing ranking order:

- 30 a) *Absent* – an assertion not satisfied by any attribute value. The attribute is absent from the presented credentials. In one embodiment said assertion type typically is used to require the absence of an entire set of credentials (e.g. “no SSL client certificate”).

- d) *Single-value* – an assertion satisfied by a single attribute value (e.g. "hair color is blue").
- e) *Multi-value* – an assertion satisfied by any value within a set of attribute values (e.g. "port number in the range of 3200 to 4200").
- 5 d) *Present* – an assertion satisfied by any attribute value, wherein the associated attribute must be present in associated presented credentials.
- e) *Ignore* – an assertion always satisfied, irrespective of whether the associated attribute is present or absent from associated presented credentials. This is a "don't care" matching rule.

10

Table K herein below shows the preferred embodiment assertion types for all operations that operate on attributes to build assertions. In the preferred embodiment when a credential specification does not include any assertions about a particular attribute then the assertion type for that attribute is ignore.

15

Table K

Operation	Assertion Type
<i>absent</i>	Absent
<i>eq</i>	Single-value
<i>ge</i>	Multi-value
<i>gt</i>	Multi-value
<i>has</i>	Multi-value
<i>ip-mask</i>	Multi-value
<i>ip-range</i>	Multi-value
<i>le</i>	Multi-value
<i>lt</i>	Multi-value
<i>member</i>	If the union has a single terminal member, the assertion type is single-value, otherwise it is multi-value
<i>prefix</i>	Multi-value
<i>present</i>	Present
<i>range</i>	Multi-value
<i>root</i>	Multi-value
<i>substring</i>	Multi-value

- In the preferred embodiment assertions in a credential specification often are combined using logical operators and, or and not. For example,

```

        ( credential Credentials_Example_1
          ( assertion
            ( and
5             ( eq ip-address 207.5.63.8 )
              ( or
                ( eq ip-port 80 )
                ( gt ip-port 443 )
              )
10          )
        )
    )
  )

```

15 In the preferred embodiment a weight assigned to a credential specification is derived from a combined weight of all assertions the credential specification comprises. An algorithm herein below is used recursively to compute a combined weight of a set of assertions operated on by a logical operator:

A. An operator `not` does not affect the weight of its operand.

20 B. An operator `and` creates a union of weights of all its operands. The weights are sorted in decreasing order of attribute rank. If multiple assertions are made about a particular attribute, use a weight of a most specific assertion and discard all other weights for that attribute. If multiple **distinct**

25 **assertions** (i.e. not identical or equivalent) are made about a particular attribute at a **same level of specificity**, the assertions are enumerated. In general, the higher a number of distinct assertions made about an attribute the more specific is a credential specification. For example, the two assertions "hair is not black" and "hair is not brown" when combined in a union

30 are more specific than either individual assertion.

C. An operator `or` results in a selection of an operand with a lowest weight. In addition said combined weight is penalized, such that it weighs less than the associated assertion with the lowest weight. If two or more assertions of

35 equal weight are combined with `or`, the combined weight is lower than that of either or any individual assertion. The rationale behind the penalty is that, in

general, combined assertions yield a larger set of principals (*i.e.* is less specific) than each assertion by itself. The weight penalty is associated with the entire credential specification, not with an individual assertion or set of assertions. Thus, for every instance of the operator `or` in the credential
 5 specification, the weight penalty is incremented by one.

In the preferred embodiment a 3-tuple represents a weight of all attribute-value assertions about a specific attribute within a credential specification. Elements in the 3-tuple are:

10

- Attribute rank
- Assertion type rank
- Attribute assertion count

15 In the preferred embodiment the 3-tuple is represented by a weight S-expression in the annotated specification language. A syntax of this expression is:

20

```
( weight <attribute> <assertion-type> <assertion-count> )
```

In the preferred embodiment ranking of assertion types is fixed and defined by the Table L following:

25

Table L

Assertion Type	Rank
absent	4
single-value	3
multi-value	2
present	1
ignore	0

In the preferred embodiment ranking of an attribute is configurable by a security administrator and must be defined prior to a compilation of a policy

specification. Attribute ranking is communicated to the policy compiler in a variety of ways. Table M herein below shows a preferred embodiment of proposed rankings for attributes used in credentials for all supported protocols. Said rankings are assumed in examples used throughout the remainder of this document. It is noted that a credential attribute agent-attributes cannot be used in a specification of an initiator or target credential and therefore need not be ranked. It is further noted that the special assertions *true* and *false*, which are allowed by the policy specification language's grammar in the preferred embodiment, do not apply to any specific attribute and, thus, are assigned a special weight consisting of a zero valued attribute rank, a zero valued assertion type rank and a zero valued attribute assertion count.

Table M

Protocol/Action	Attribute	Rank
IP/ASSOCIATION UDP/ASSOCIATION ICMP/ASSOCIATION TCP/CONNECT	mac-address	3
	ip-address	2
	ip-port	1
SSL/HANDSHAKE	der-cert	5
	x509-subject	4
	x509-issuer	3
	x509-cert-path	2
SSL/HANDSHAKE	cert-status	1
HTTP/GET HTTP/POST HTTP/HEAD	http-username	3
	http-password	2
	url	1

In the preferred embodiment an attribute assertion count starts at zero for a first assertion and is incremented monotonically for all subsequent assertions. That is, the count enumerates additional assertions for the attribute. In the preferred embodiment the assertion count is omitted from the weight S-expression when said count is zero.

In the preferred embodiment a weight S-expression is omitted when an assertion type is ignore.

In the preferred embodiment the three elements of a 3-tuple are used in sorting a collection of 3-tuples. The attribute rank as a primary key, the assertion type rank as a secondary key, and the attribute assertion count as a tertiary key produce an ordered list of 3-tuples sorted in decreasing order of rank and count. In the preferred embodiment said sorted list is used to rank credential specifications against each other. The sorting algorithm is described using pseudo-code in Table N herein below:

Table N

```

20 Sort_3tuples: subroutine ( 3tuple A, 3tuple B )
    begin
        if A.attribute_rank > B.attribute_rank
            return (A is greater than B);
25     else if A.attribute_rank < B.attribute_rank
            return (A is less than B);
        else // same attribute rank
            if A.assertion_type > B.assertion_type
                return (A is greater than B);
30     else if A.assertion_type < B.assertion_type
            return (A is less than B);
        else // same assertion type
            if A.assertion_count > B.assertion_count
                return (A is greater than B);
35     else if A.assertion_count < B.assertion_count
            return (A is less than B);
        else // same assertion count
            return (A is equal to B);

```


end

A weight penalty is represented by the following S-expression in the annotated specification language:

5

```
( weight-penalty <penalty-count> )
```

where <penalty-count> is an integer representing a number of or operators in a credential specification.

10

Thus, Credentials_Example_1 herein above is annotated as follows:

```
( weight ip-address single-value )  
( weight ip-port multi-value )  
15 ( weight-penalty 1 )
```

In the preferred embodiment a credential specification can combine previous credential specifications with each other or with additional assertions. In the preferred embodiment rules for a combination of assertions with logical operators apply equally to a combination of credential specifications. For example:

```
( credential Credentials_Example_2  
  ( assertion  
25    ( eq ip-address 207.5.63.22 )  
  )  
)
```

```
( credential Credentials_Example_3  
  ( assertion  
30    ( and  
      Credentials_Example_2  
      ( gt ip-port 42 )  
    )  
  )  
35 )  
)
```

```
( credential Credentials_Example_4
```

```

      ( assertion
        ( and
          ( or
            Credentials_Example_1
5          Credentials_Example_3
          )
          ( lt ip-port 1025 )
        )
      )
10    )

```

The weight of Credentials_Example_2 is:

```
( weight ip-address single-value )
```

15 The weight of Credentials_Example_3 is:

```
( weight ip-address single-value )
```

```
( weight ip-port multi-value )
```

In the embodiment to compute the weight of Credentials_Example_4 first
 20 compute a weight of the or expression. Credentials_Example_1 is
 selected as having a lowest weight because of an associated weight penalty.
 Furthermore, the or expression in Credentials_Example_4 increases the
 weight penalty further, yielding:

```

      ( weight ip-address single-value )
25    ( weight ip-port multi-value )
      ( weight-penalty 2 )

```

In the embodiment the and expression adds an additional, distinct, assertion
 about ip-port. The assertion is of the same type as one currently selected
 30 because they are both multi-value assertions. The assertion count for ip-
 port is incremented, yielding:

```

      ( weight ip-address single-value )
      ( weight ip-port multi-value 1 )
      ( weight-penalty 2 )
35

```

In the embodiment a ranking algorithm for comparing and ordering credentials is implied in the example previously described herein above. Following in Table O is an associated algorithm using pseudo-code:

Table 0

```

Rank_credentials: subroutine ( credentials A, credentials B )
begin
    Let X be a sorted list of 3-tuples that describe the weight of
    credential specification A;

    Let Y be a sorted list of 3-tuples that describe the weight of
    credential specification B;

    // compare 3-tuples

    for all i in X do
        R = Sort_3tuples(X[i],Y[i]); // defined above
        if R is greater-than
            return (A ranks higher than B);
        else if R is less-than
            return (A ranks lower than B);
        else
            continue; // 3-tuples are equal
    end

    // X and Y are the same; compare weight penalties

    if A.weight_penalty < B.weight_penalty
        return (A ranks higher than B);
    else if A.weight_penalty > B.weight_penalty
        return (A ranks lower than B);
    else
        return (A and B have the same rank);
end

```

The following Table P ranks example credentials according to the preferred embodiment using the algorithm herein above. A weight column shows 3-tuples using a format $W:x,y,z$, wherein x is an integer value for an attribute

- rank (Table M), y is an integer value for an assertion type (Table P), and z is an assertion count. A weight penalty is shown as $P:x$, wherein x is a penalty count. It is noted that the higher a rank of a credential specification, the more specific it is. For completeness, the table includes ranking for built-in
- 5 credentials denoted by absent, present and ignore. Said built-in credentials make assertions about and in the order of an absence, presence, and irrelevance of any credentials presented by a protocol event. It is noted that in the preferred embodiment ignore and present always rank lower and absent higher than do any user-defined credentials.

10

Table P

Name	Weight	Rank
absent (built-in)	W:*,5	6
Credentials_Example_4	W:2,3 W:1,2,1 P:2	5
Credentials_Example_3	W:2,3 W:1,2	4
Credentials_Example_1	W:2,3 W:1,2 P:1	3
Credentials_Example_2	W:2,3	2
present (built-in)	W:*,1	1
ignore (built-in)	W:*,0	0

Ordering of Rules

- 15 In the preferred embodiment policy rules must be organized such that when two or more rules are satisfied by a protocol event, the most specific rule for that event is selected. The specificity of a policy rule is fully determined by the specificity of the credentials it uses.
- 20 In the preferred embodiment policy rules are organized as follows:

- 1) Rules are segregated by protocol. For example, rules that apply to a TCP protocol are separated from those that apply to a SSL protocol.
- 2) Within each protocol group, rules are segregated by action. For example,
 5 rules that only apply to a TCP CONNECT action are separated from those that only apply to a TCP CLOSE action.
- 3) Within each protocol-action group, rules are ranked by the specificity of their respective credentials. The ranking algorithm is:
 10
 - a) Create a 2-tuple from a ranking order of an initiator credential and a target credential. The first element in said 2-tuple is a highest-ranking value and the second element a lowest. That is, said 2-tuple is defined as $(\text{MAX}(I,T), \text{MIN}(I,T))$, wherein I and T are the ranking values for the
 15 initiator and target credentials, respectively.
 - b) Sort the rules in increasing ranking order using the first element in the 2-tuple as the primary sorting key and the second element as the secondary key. Rules with identical 2-tuples are given the same ranking
 20 number. The rule or rules with the highest-ranking number is the most specific rule for the protocol group.

In the preferred embodiment and because rules are ranked directly from the ranking of their credentials, a special representation is not provided in the
 25 annotated specification language for the ranking of the policy rules.

Following is an example using credentials from herein above:

```

30  ( rule Rule_Example_1
      ( protocol TCP )
      ( action CONNECT )
      ( initiator Credentials_Example_2 ) // ranked #2
      ( target Credentials_Example_1 )   // ranked #3
      ( outcome
35      ...
      )
  )

```

```

    )

    ( rule Rule_Example_2
      ( protocol TCP )
      ( action CONNECT )
      ( initiator Credentials_Example_1 ) // ranked #3
      ( target Credentials_Example_2 )   // ranked #2
      ( outcome
        ...
      )
    )

    ( rule Rule_Example_3
      ( protocol TCP )
      ( action CONNECT )
      ( initiator Credentials_Example_2 ) // ranked #2
      ( target Credentials_Example_4 )   // ranked #5
      ( outcome
        ...
      )
    )
  
```

Table Q herein below shows how said rules are ranked according to the invention.

Table Q

Name	Credentials Rank	Rule Rank
Rule_Example_3	(T:5, I:2)	2
Rule_Example_1	(T:3, I:2)	1
Rule_Example_2	(I:3, T:2)	1

It is noted that Rule_Example_1 and Rule_Example_2 are ranked at the same specificity level. This does not represent a problem because the respective initiator and target credential sets are non-intersecting and used in different roles.

In the preferred embodiment it is possible for two or more rules at a same specificity level to be satisfied by a single protocol event. During policy specification a security administrator disambiguates the evaluation of rules with the same specificity level by forcing a ranking order among them.

- 5 Forcing a ranking order is done by specifying that one rule is ranked above another rule and is termed forced ranking. Forced ranking is expressed by means of the following S-expression:

```
( rank-above <rule-name> )
```

- 10 For example, to give Rule_Example_2 precedence over Rule_Example_1, the following S-expression is added to a definition of Rule_Example_2:

```
( ( rank-above Rule_Example_1 )
```

- In the preferred embodiment after performing the standard ranking algorithm
15 herein above, the Policy Engine 101 evaluates all rank-above expressions and reassigns ranking numbers to each rule accordingly. In the preferred embodiment it is important to note that forced ranking does not force a ranking of an affected rule to a level of a more specific rule higher in the ranking order. Instead a new ranking level is created for the affected rule and
20 all other ranking numbers of more specific rules are incremented accordingly.

For example, Rule_Example_2 herein above is given ranking number 2 and the ranking number of Rule_Example_3 herein above is incremented from 2 to 3.

- 25 In the preferred embodiment forced ranking is applied to any rule and is not limited by rules having only non-unique ranking numbers. In this embodiment security administrators are cautioned not to use said forced ranking feature unless absolutely necessary. Its misuse may result in a policy specification that is both difficult to manage and difficult to evaluate. In the preferred
30 embodiment runtime conflicts in the evaluation of rules (i.e. when a protocol event is satisfied by multiple rules) typically can be solved by redesigning credentials upon which said rules are based. Useful tips are provided herein below.

Evaluation Algorithm

In the preferred embodiment the Policy Engine 101 applies a policy evaluation algorithm to each incoming protocol event. The algorithm results in a selection of a policy rule applicable to the protocol event and may produce an immediate or final disposition.

Following is a step-by-step description of the evaluation algorithm according to the preferred embodiment. It is noted that the evaluation procedure described herein below is in conceptual form and does not take into account any possible runtime optimizations:

- 1) Select a set of rules applicable to an Agent reporting an event;
- 2) From said set, select a second set of rules applicable to an associated examined protocol.
- 3) From said second set, select a third set of rules applicable to an associated examined protocol action.
- 4) Starting with a most specific policy rule in said third set and descending to a least specific rule find a policy rule satisfied by said protocol event. A matching algorithm according to the preferred embodiment is as follows:
 - a) If one or more orderly listed prerequisite rules are specified, ensure at least one of said prerequisite rules is satisfied by a previously processed protocol event. In the preferred embodiment a prerequisite rule is satisfied if it is a pending policy rule for the protocol event.
 - b) Match initiator and target credentials in the policy rule against the corresponding initiator and target credentials presented in the protocol event.

- 5) If a policy rule satisfying the protocol event is not found the Policy Engine 101 generates a disposition for the network event indicating that a policy specification error was encountered. Effectively the processing of the network event thereby terminates.
- 5
- 6) If a policy rule satisfying the protocol event is found, the Policy Engine 101 checks for other rules having a same ranking number and also satisfying the event. If such rules are found the Policy Engine 101 uses the following algorithm in the preferred embodiment to select a single applicable rule:
- 10
- a) Rules that specify all protocols (*i.e.* using ignore or present) are less specific than rules that explicitly list a set of one or more protocols.
- b) Rules that specify all actions (*i.e.* using ignore or present) are less specific than rules that explicitly list a set of one or more actions.
- 15
- c) Rules that have prerequisites are more specific than rules that do not have prerequisites. Rules that specify a higher-ranking prerequisite are more specific than rules that specify a lower-ranking prerequisite. In the preferred embodiment a ranking relationship is relevant only if both prerequisite rules belong to a same protocol-action group.
- 20
- d) If thereafter a single rule is determined as more specific than the others it is selected for the protocol event. If more than one rule remains the Policy Engine 101 sorts the remaining rules in increasing lexical order by name and selects a first rule from the sorted rules having an immediate disposition indicating in decreasing order of precedence:
- 25
- i) a policy violation (any disposition code other than OK or CONTINUE);
- 30
- ii) CONTINUE (allows other rules to examine further the network event); and
- iii) OK

The outcome of the policy evaluation algorithm herein above is a policy rule that satisfies the protocol event. If an immediate outcome is specified for that rule, it is executed, producing a disposition for the protocol event. If the disposition comprises a final disposition code (any code other than
5 CONTINUE), the disposition is also the final disposition for the network event.

Otherwise in the preferred embodiment the selected policy rule is a pending policy rule for the network event. In absence of any further protocol events the pending policy rule is promoted to selected policy rule. A final outcome of
10 the selected policy rule is executed producing a final disposition for the network event.

Policy Specification Guidelines

Provided herein below in Table R are a number of practical guidelines coupled to the preferred embodiment for the development and specification phases of a security policy. Adhering to the guidelines ensures efficient and accurate evaluation of a policy by the Policy Engine 101. It is intended to incorporate the guidelines into a graphical policy editing invention using wizards, policy templates and other UI mechanisms that among other uses
20 simplify and direct the policy specification process.

Table R

Rule #1: Work on group relationships

25 The first step in policy specification is identifying the communicating entities and resources that interact with each other over the network, that is to say, specifying the credentials for both initiator and target principals. Defining groups in relation to each other can significantly enhance the ranking of credentials. This is best done by:

- 30
- Defining all large groups first (e.g. all hosts in the corporate network, all valid certificates).
 - Defining all other groups by subsetting larger groups (e.g. all hosts in the marketing subnetwork, all certificates issued by the corporate CA, all revoked certificates).
- 35 The process of defining a group as a subset of another can be thought of as the process of specializing the credentials specification for the larger group. Thus, the smaller group's credentials are more

specific than those of the larger group. Likewise, creating a larger group through the union of smaller groups generalizes the credentials specification of the smaller groups, thus resulting in less specific credentials for the larger group.

5 Rule #2: Deny first, allow later

A good security management principle is that of denying access to a resource unless access is explicitly granted. Thus, when specifying a network's security policy the first step must be to deny access to all target principals via rules that identify initiators via the broadest possible credentials. One can then grant access to each target principal solely to the group of principals to which access should be granted.

10

For example, to protect a set of host machines from access by all but a small set of principals, one can define a rule that denies access to these machines and whose initiator is denoted by ignore. A second rule allowing access can then be defined. It specifies the same target principal and, as the initiator, a credential specification that describes, in the narrowest possible manner, the principals being granted access. The ranking algorithm guarantees that the rule granting access ranks higher than the rule denying it.

15

20

It is crucial that the credential specification for the principals being granted the access privilege be as specific as possible, with all other principals being denied access. This ensures that access is not inadvertently granted to non-privileged principals.

25

In general, the first policy rule in every protocol layer is one that denies access to all and by all communicating entities (using ignore for both initiator and target principals) for all protocol actions (again using ignore).

30

Rule #3: Prerequisites are your friends, use them often and use them wisely

Prerequisite rules can play a critical role in the disambiguation of like-ranked rules. Thus, prerequisites should be used whenever possible. In particular, prerequisites should be used in a way that targets each rule to the smallest set of principals possible, and that prevents the repetition of credentials within a set of related rules. For example, if an IP rule exists that defines communication between hosts in two subnets and we want to define a TCP rule affecting the same set of hosts, we should define a TCP rule that takes the aforementioned IP rule as a prerequisite. In addition, the credentials used in the TCP rule should not include assertions that repeat what has already been established by the IP rule (e.g. the IP addresses of the relevant hosts). Instead the TCP rule credentials should specialize (if so desired) the specification of the host credentials, e.g. limiting the host services covered by the rule (i.e. stating the IP ports of interest).

35

40

Rule #4: Make dispositions final, unless they are not

Immediate outcomes that produce a final disposition should be used whenever possible. In other words, unless one knows that a rule at a given protocol layer may be overridden by a specific rule at a higher protocol layer, the immediate outcome for the former rule should always produce the final

disposition for the network event. This prevents a rule's outcome from being inadvertently subsumed by another protocol event.

- 5 In general, unless a rule is explicitly listed as a prerequisite rule for another rule higher in the protocol stack, its immediate outcome should produce the final disposition for the network event.

Rule #5: If you know the Agent, name it

If a policy rule only applies to communications within a specific network segment, restrict the rule's scope by specifying the Agent(s) reporting the protocol events for which this rule should be considered.

10

By doing so, one eliminates that rule from being considered in events reported by other Agents.

- Although the invention is described herein with reference to a variety of preferred embodiments, one skilled in the art will readily appreciate that other applications may be substituted for those set forth herein without departing from the spirit and scope of the present invention. Accordingly, the invention should only be limited by the Claims included below.
- 15

CLAIMS

1. A declarative language system for specifying in an annotated policy
5 specification a security policy of a network event, wherein said network event
comprises a stack having a plurality of protocol events, wherein each of said
plurality of protocol events is associated with a predefined protocol layer, and
wherein said network event is an interaction between an active principal and a
passive principal, said declarative language system comprising:
- 10 a declarative language comprising a plurality of objects, such that each
object of said plurality of objects comprises at least one list having a first
element;
a declarative language editor for providing means for specifying in a
first
15 policy specification said security policy using said declarative language;
a declarative language compiler for providing means for compiling said
first policy specification and generating said annotated policy specification;
means for loading said annotated policy specification into a Policy
Engine;
20 means for said Policy Engine to receive said network event from an
Agent;
means for said Policy Engine to evaluate said security policy against
said
network event and to generate a disposition for said network event;
25 means for said Policy Engine to communicate agent directives to said
Agent; and
means for said Policy Engine to output said network event and said
disposition to a datastore.
- 30 2. The system of Claim 1, wherein said declarative language uses the S-
expression language.
3. The system of Claim 2, wherein the S-expression language is a variant
by Rivest used in SPKI/SDSI.

4. The system of Claim 3, wherein for said each object, said first element of said list is a type, such that said type is associated with said each object.
5. The system of Claim 4, wherein said type is a byte string.
- 5 6. The system of Claim 3, wherein a canonical representation of said S-expression language is supported.
7. The system of Claim 3, wherein an advanced representation of said S-expression language is supported.
- 10 8. The system of Claim 6, wherein said canonical representation is digitally signed.
- 15 9. The system of Claim 2, wherein said declarative language allows embedded comments in said S-expression language.
10. The system of Claim 2, wherein said declarative language supports macros.
- 20 11. The system of Claim 2, wherein said declarative language supports included files.
12. The system of Claim 1, wherein said each object is a first-class object.
- 25 13. The system of Claim 12, wherein said first-class object is a built-in object, such that said built-in first-class object is associated with said declarative language compiler, and wherein said built-in first-class object is unextendable within said annotated policy specification.
- 30 14. The system of Claim 12, wherein said first-class object is a user-defined object.
15. The system of Claim 12, wherein said first-class object further comprises a description field.
- 35

16. The system of Claim 12, wherein said first-class object is any of:
a policy;
a group;
5 a credential, said credential having a specificity;
a condition;
a disposition; and
a rule, said rule having an outcome.
- 10 17. The declarative language of Claim 16, wherein said policy comprises:
a name parameter for referencing said security policy;
a version number parameter associated with a version of said
declarative language; and
at least one rule of a plurality of rules.
- 15 18. The system of Claim 17, wherein said group comprises:
a union object, said union object comprising a plurality of items, each
item having a same type;
a name parameter for said union object; and
20 a type parameter for defining said same type of said plurality of items.
19. The system of Claim 18, wherein one of said plurality of items is a
second group comprising a second union object, said second union object
comprising a second plurality of items, each item of said second plurality of
25 items having said same type.
20. The system of Claim 18, wherein said type is a primitive data type in
said declarative language.
- 30 21. The system of Claim 20, wherein said primitive data type includes, but
is not limited to, any of:
a string;
an IP address;
a MAC address;
35 an integer;

a version number; and
a hash value.

22. The system of Claim 16, wherein said credential comprises:
5 a name parameter for referencing said credential ; and
at least one assertion, wherein said assertion is a logical expression
comprising:
a plurality of attributes, each of said plurality of attributes having an
attribute-value; and
10 a plurality of logical operands;
wherein said credential is associated with one of said active principal
and said passive principal in said event, and wherein said credential has a
specificity; and
wherein said credential is associated with a first protocol, said first
15 protocol having a set of associated attributes and a set of associated
operands.
23. The system of Claim 22, wherein each of plurality of credential
attributes comprises an implied fetching function, said function returning a
20 value associated with each of said plurality of credential attributes.
24. The system of Claim 23, wherein said function is argumentless and
wherein said returned value is a single value.
25. The system of Claim 23, wherein said function has a plurality of
25 arguments and wherein said returned value is a union of a plurality of values.
26. The system of Claim 22, wherein said credential further comprises a
second credential, wherein said second credential is associated with said first
30 protocol.
27. The system of Claim 16, wherein said condition comprises:
a name parameter for referencing said condition; and
at least one assertion, wherein said assertion is a logical expression
35 comprising;

a plurality of attributes, each of said plurality of attributes having an attribute value;

a plurality of logical operands;

wherein said condition defines a constraint on said event; and

- 5 wherein said condition is associated with a first protocol, said first protocol having a first set of associated attributes and a first set of associated operands.

28. The system of Claim 27, wherein said condition further comprises a
10 second condition wherein said second condition is associated with said first protocol.

29. The system of Claim 16, wherein said disposition comprises:
a disposition code for indicating one of the absence of a violation of
15 said rule and the presence of said violation of said rule; and
wherein said disposition represents said outcome of said rule.

30. The system of Claim 29, further comprising:
a logging directive, wherein said logging directive comprises:
20 a severity code, said severity code indicating a severity level of said disposition; and
a human readable string for providing additional details.

31. The system of Claim 30, wherein said severity code is used by a
25 logging subsystem for classifying and filtering said network event.

32. The system of Claim 29, further comprising an agent directive having instructions that are communicated to said Agent.

- 30 33. The system of Claim 32, wherein said Agent monitors network traffic.

34. The system of Claim 32, wherein said Agent enforces security policy.

- 35 35. The system of Claim 16, wherein said rule for evaluating said event comprises:

- a protocol field associated with said event;
a plurality of actions associated with said event;
an initiator for representing said active principal of said event;
a target for representing said passive principal of said event, and
5 means for said outcome to generate a disposition by specifying
constraints upon said event, said outcome comprising:
at least one of a plurality of conditional statements and a default
statement, wherein each of said plurality of conditional statement comprises a
keyword and a disposition, and wherein said plurality of conditional
10 statements are evaluated in chronological order.
36. The system of Claim 35, further comprising:
an agent field for representing said Agent associated with said event,
wherein said agent field is associated with an agent credential and wherein
15 said rule is applied when said agent credential is satisfied.
37. The system of Claim 35, further comprising:
a prerequisite having a plurality of rules, such that said prerequisite is
satisfied when at least one of said plurality of rules is applied to a prior event.
20
38. The system of Claim 35, wherein said outcome comprises any of an
immediate outcome and a final outcome, wherein said immediate outcome is
evaluated by said Policy Engine when said rule is selected, and wherein said
final outcome is evaluated when said Policy Engine determines said event is
25 final.
39. The system of Claim 1, further comprising:
an annotated specification language;
wherein said first policy specification further comprises:
30 a plurality of credentials,
a plurality of conditions,
a plurality of rules;
wherein means for compiling comprises:
means for checking said first policy specification for syntax errors and
35 semantics errors;

means for checking said first policy specification for credential errors;
means for checking said first policy specification for condition errors;
means for checking said first policy specification for completeness and
coverage of said plurality of rules;

5 means for ordering said plurality of credentials by using said annotated
specification language, whereby for each of said plurality of credentials a
credential rank is determined; and

means for ordering said plurality of rules by using said annotated
specification language.

10

40. The system of Claim 39, further comprising:
an annotated specification language for providing additional information
to
said means to evaluate;

15 means for Policy Engine to receive said plurality of protocol events and
to
provide a sequencing of said plurality of protocol events by using said
associated predefined protocol layers;

means for Policy Engine to select a policy rule associated with each of
20 said plurality of protocol events, using a specificity of said policy rule;
means for Policy Engine to determine said policy rule outcome;
means for Policy Engine to render said policy rule as a pending policy
rule;

25 means for Policy Engine to render one of said policy rule and said
pending policy rule as final.

41. The system of Claim 40, wherein said Policy Engine is adapted to
receive and evaluate incomplete data in any of said plurality of protocol
events.

30

42. The system of Claim 39, wherein said means for ordering said plurality
of credentials further comprises:

means for computing a combined weight for each of said plurality of

credentials of each attribute weight, having a plurality of attribute-value assertions of said plurality of credential attributes, wherein each attribute weight comprises:

- an attribute rank;
- 5 an assertion type rank;
- an attribute assertion count;
- means for computing a second combined weight of a subset of said plurality of attribute-value assertions operated on by a logical operator;
- means for computing a credential weight penalty for each of said
- 10 plurality of credentials; and
- means for comparing said plurality of credentials.

43. The system of Claim 42, wherein said attribute weight is represented by a 3-tuple having a weight keyword in said annotated specification
15 language.

44. The system of Claim 42, wherein said logical operator is any of and, or, and not.

20 45. The system of Claim 42, wherein means for compiling comprises means for a security administrator to configure said attribute rank.

46. The system of Claim 42, wherein said attribute assertion count starts at zero and is incremented monotonically for subsequent assertions.

25

47. The system of Claim 42, wherein said assertion count is zero, said attribute assertion count is omitted from said 3-tuple.

48. The system of Claim 43, further comprising means to sort a plurality of
30 3-tuples, wherein said attribute rank is a primary key, said assertion type rank is a secondary key, and said attribute assertion count is a tertiary key, thereby providing a sorted list.

49. The system of Claim 42, wherein means for computing a weight
35 penalty comprises:

a weight-penalty keyword in said annotated specification language having a penalty-count parameter, wherein said penalty-count parameter is an integer representing a total number of occurrences of logical operator or in each credential.

5

50. The system of Claim 42, wherein means for comparing said plurality of credentials comprises:

means to determine a highest ranking 3-tuple from said sorted 3-tuples;

means to compare credential weight penalties; and

10 means to assign said credential rank.

51. The system of Claim 39, wherein means for ordering said plurality of rules comprises:

a plurality of predetermined protocols;

15 a plurality of predetermined protocol-action groups;

means to assign each of said rules to one of said predetermined protocols;

means to assign each of said rules to one of said predetermined protocol-action groups;

20 means to rank each of said rules in said predetermined protocol-action groups by using said credential ranking value for said target credential of said rule and by using said credential ranking value for said initiator credential of said rule;

25 means to sort in increasing order each of said ranked rules in said predetermined protocol-action groups.

52. The system of Claim 51, further comprising:

means to force said rule ranking value for any of each of said rules using

30 said annotated specification language, said annotated specification language having a rank-above expression having a rule-name parameter.

53. The system of Claim 51, further comprising:

a 2-tuple for each said rule, said 2-tuple having a first element and a

35 second element;

wherein said first element is a highest credential ranking value of said target credential and initiator credential; and

wherein said second element is a lowest credential ranking value of said target credential and initiator credential; and

5 wherein means to sort uses said 2-tuple of each rule.

54. The system of Claim 52, wherein means to force said rule ranking value comprises:

generating a new ranking level for said forced ranked rule, whereby
10 each of said rules having a rule ranking level at forced level or higher are incremented.

55. A method for evaluating a policy using a plurality of policy rules, each rule having a ranking and a disposition, to a protocol event reported by an Agent, said protocol event having a protocol, a protocol action, a target
15 credential, and an initiator credential, comprising the steps of:

selecting a first set of rules from said plurality of policy rules, such that each rule is associated with said Agent;

selecting a second set of rules from said first set of rules, such that
20 each rule is associated with said protocol from said event;

selecting a third set of rules from said second set of rules, such that each rule is associated with said protocol action from said event;

searching for a most specific policy rule from said third set, such that said most specific policy rule is satisfied by said protocol event and generating
25 an error disposition when said most specific policy rule is undetermined;

checking said third set of rules for a fourth set of rules having same said ranking as said selected most specific policy rule;

providing means to select a single applicable rule from said fourth set of rules.
30

56. The method of Claim 55, further comprising the step of:

producing a final disposition for a network event, wherein said network event comprises said protocol event.

57. The method of Claim 55, wherein the step of searching for a most specific policy rule further comprises the steps of:

satisfying any of a plurality of prerequisite rules by a previous protocol event in an order corresponding to an order of said plurality of prerequisite

5 rules; and

matching a rule target credential and a rule initiator credential with said event target credential and said event initiator credential.

58. The method of Claim 55, wherein the step of providing means to select
10 a single applicable rule further comprises the steps of:

designating any rule of said fourth set of rules that specifies all of said plurality of protocols as less specific;

designating any rule of said fourth set of rules that specifies all of said plurality of protocol actions as less specific;

15 designating any rule of said fourth set of rules having prerequisite rules as more specific, wherein a rule having a higher ranking prerequisite is more specific than a rule having a lower ranking prerequisite;

20 sorting any remaining rules in increasing lexical order by said names and thereafter by said immediate dispositions in decreasing order of precedence; and

selecting said single applicable rule from first rule of said sorted rules.

59. The method of Claim 58, wherein said immediate dispositions in decreasing order of precedence comprises:

25 a policy violation;

CONTINUE; and

OK.

60. A method for processing an outcome of a policy rule associated with a
30 protocol event of a network event, comprising the steps of:

if said outcome is specified and immediate, executing said outcome, producing thereby a disposition for said protocol, and designating said disposition final for said network event if said disposition comprises a final disposition code;

designating said policy rule a pending policy rule for said network event;

promoting said pending policy rule to selected policy rule, if further protocol events are absent;

- 5 executing a final outcome of said selected policy rule; and
 producing a final disposition for said network event of said selected policy rule final outcome.

61. A computer implemented system for interpreting different protocols,
10 comprising:

 a language;

 a policy editor adapted to use said language;

 a policy specification generated by said policy editor and written in said language;

- 15 a policy engine receiving at least one event, wherein said policy engine is associated with said policy specification, and wherein said policy engine interprets said language; and

 a disposition generated by said policy engine using said event and said policy specification.

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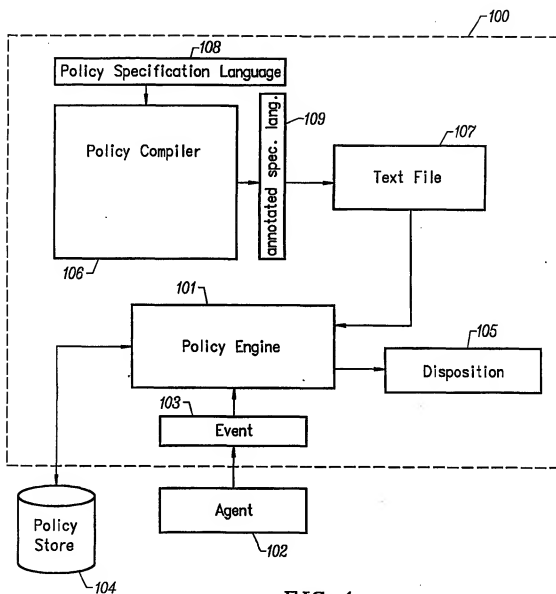


FIG. 1

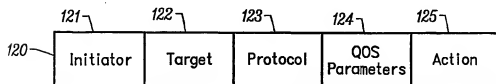


FIG. 2

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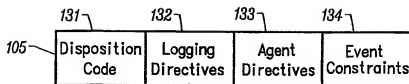


FIG. 3

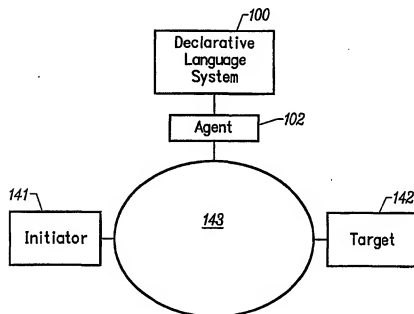


FIG. 4

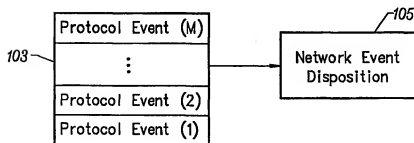


FIG. 5A

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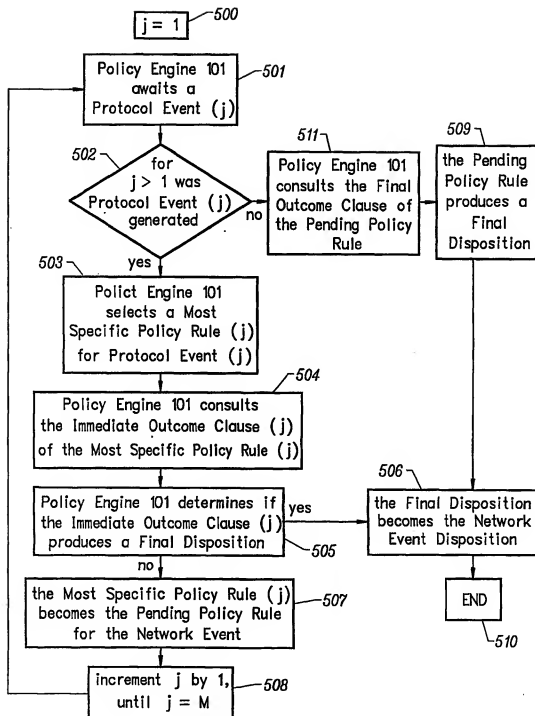


FIG. 5B

EP1024627

Publication Title:

A method and apparatus for managing a firewall

Abstract:

Abstract of EP1024627

A method and apparatus are disclosed for managing a firewall. The disclosed firewall manager facilitates the generation of a security policy for a particular network environment, and automatically generates the firewall-specific configuration files from the security policy simultaneously for multiple gateways. The security policy is separated from the vendor-specific rule syntax and semantics and from the actual network topology. Thus, the security administrator can focus on designing an appropriate policy without worrying about firewall rule complexity, rule ordering, and other low-level configuration issues. In addition, the administrator can maintain a consistent policy in the presence of intranet topology changes. The disclosed firewall manager utilizes a model definition language (MDL) and an associated parser to produce an entity relationship model. A model compiler translates the entity-relationship model into the appropriate firewall configuration files. The entity-relationship model provides a framework for representing both the firewall-independent security policy, and the network topology. The security policy is expressed in terms of "roles," which are used to define network capabilities of sending and receiving services. A role may be assumed by different hosts or host-groups in the network. A visualization and debugging tool is provided to transform the firewall-specific configuration files into a graphical representation of the current policy on the actual topology, allowing the viability of a chosen policy to be evaluated. A role-group may be closed to prevent the inheritance of roles.

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